

Let the **BEAST** of **CRIME** and **TIME**
be not so **LUCKY**

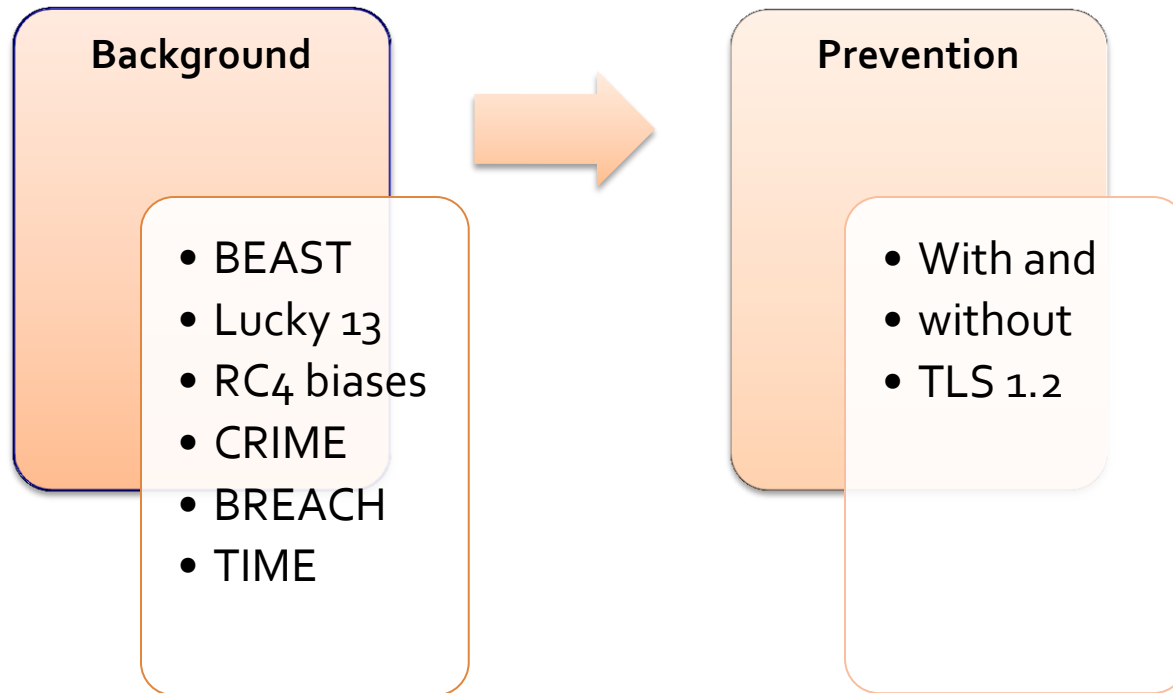


TL\$



Pratík Guha Sarkar

Takeaway



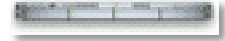
BEAST



BEAST - What is it ?

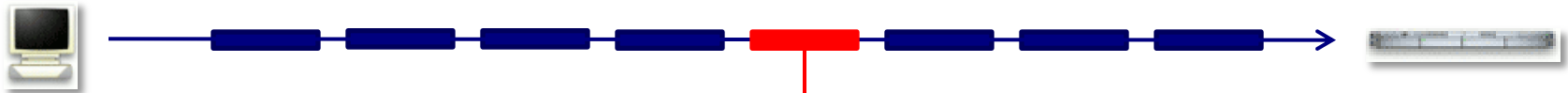
- **Browser Exploit Against SSL/TLS**
- Used crypto flaws in SSL to recover plaintext cookies
- Refined previous attacks on CBC in SSL to make them practical
 - Innovation was exploiting **chosen boundary** capability

CBC Attack in SSL



CBC Attack in SSL


Target Block For Decrypting



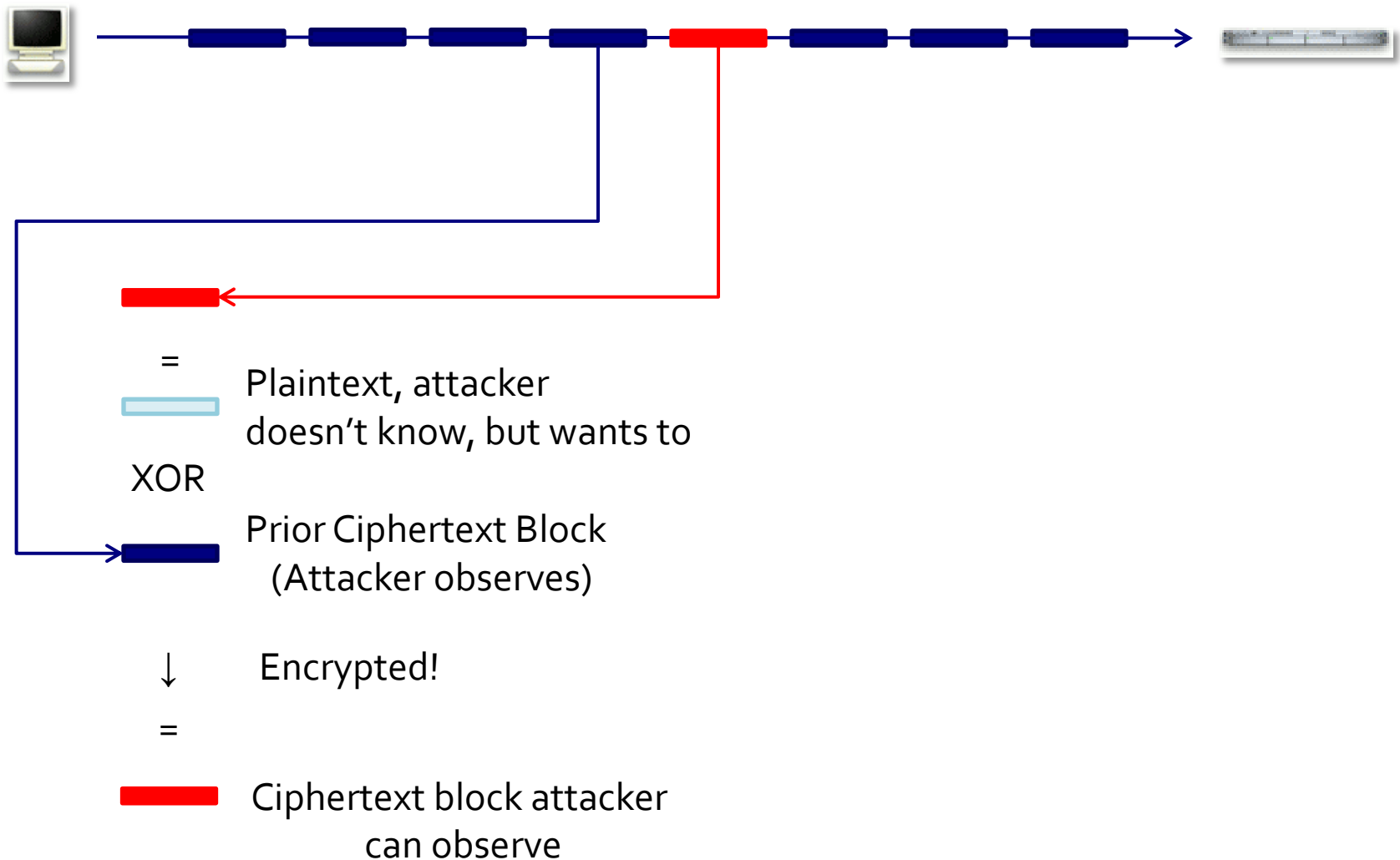
=

?

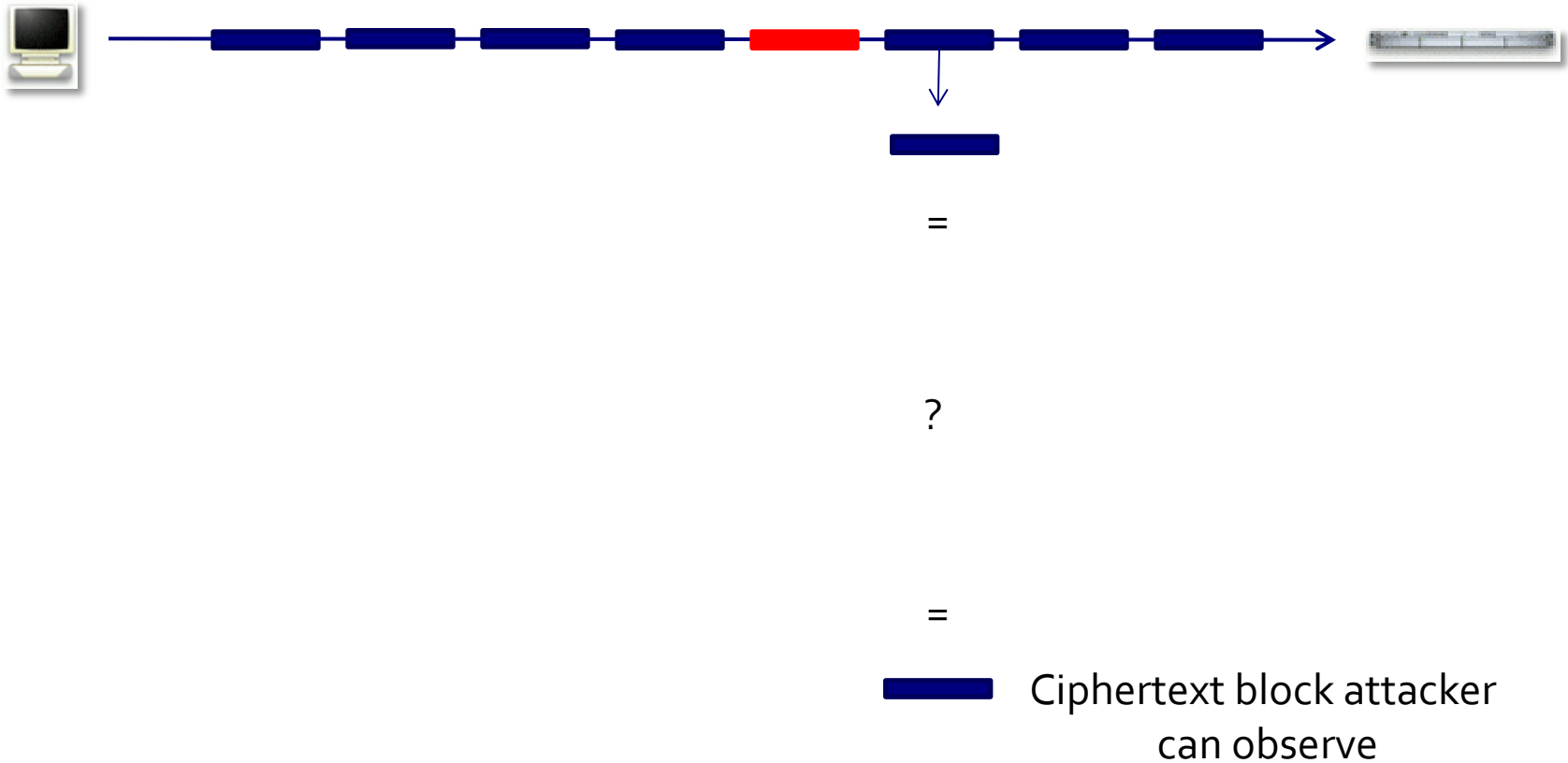
=

 Ciphertext block attacker
can observe

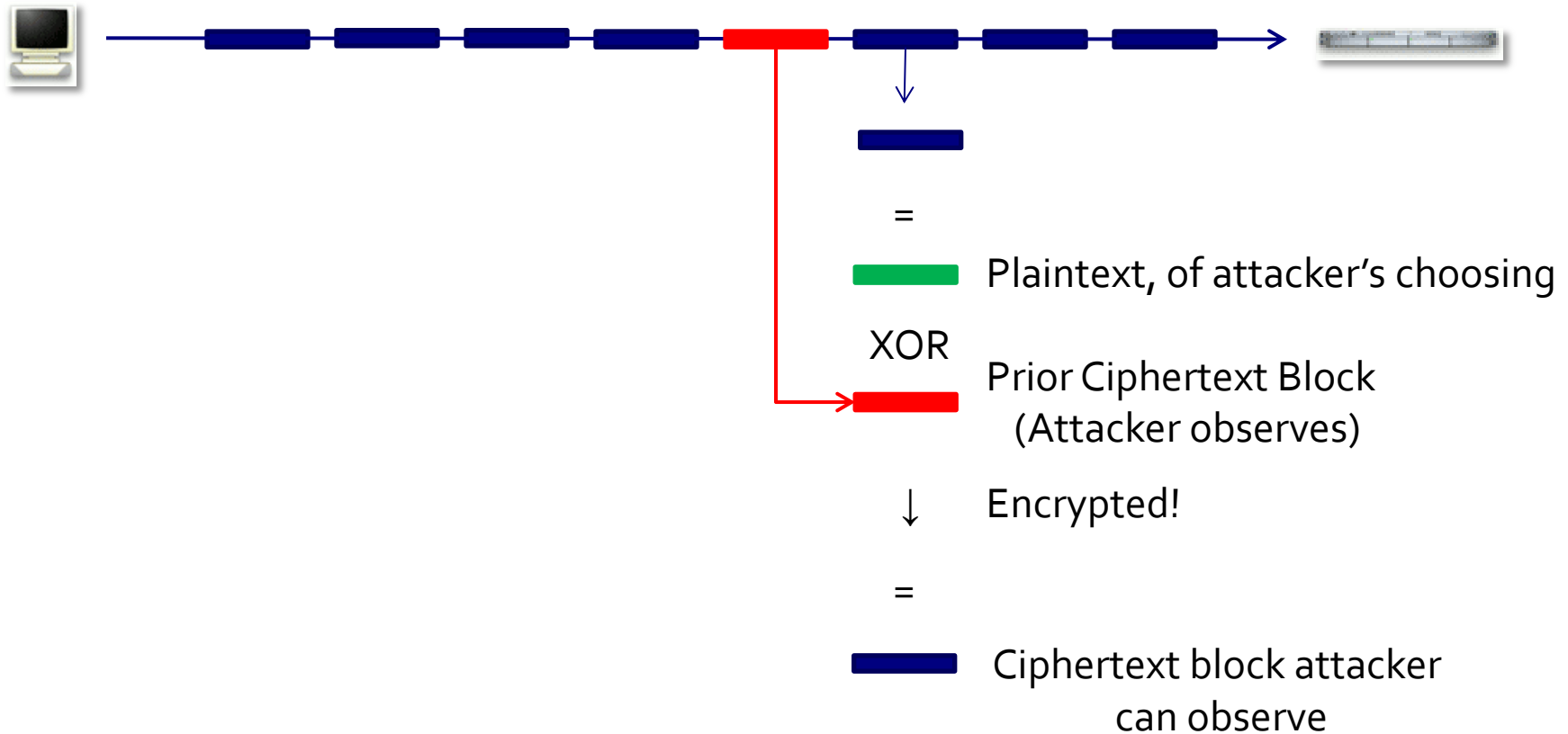
CBC Attack in SSL



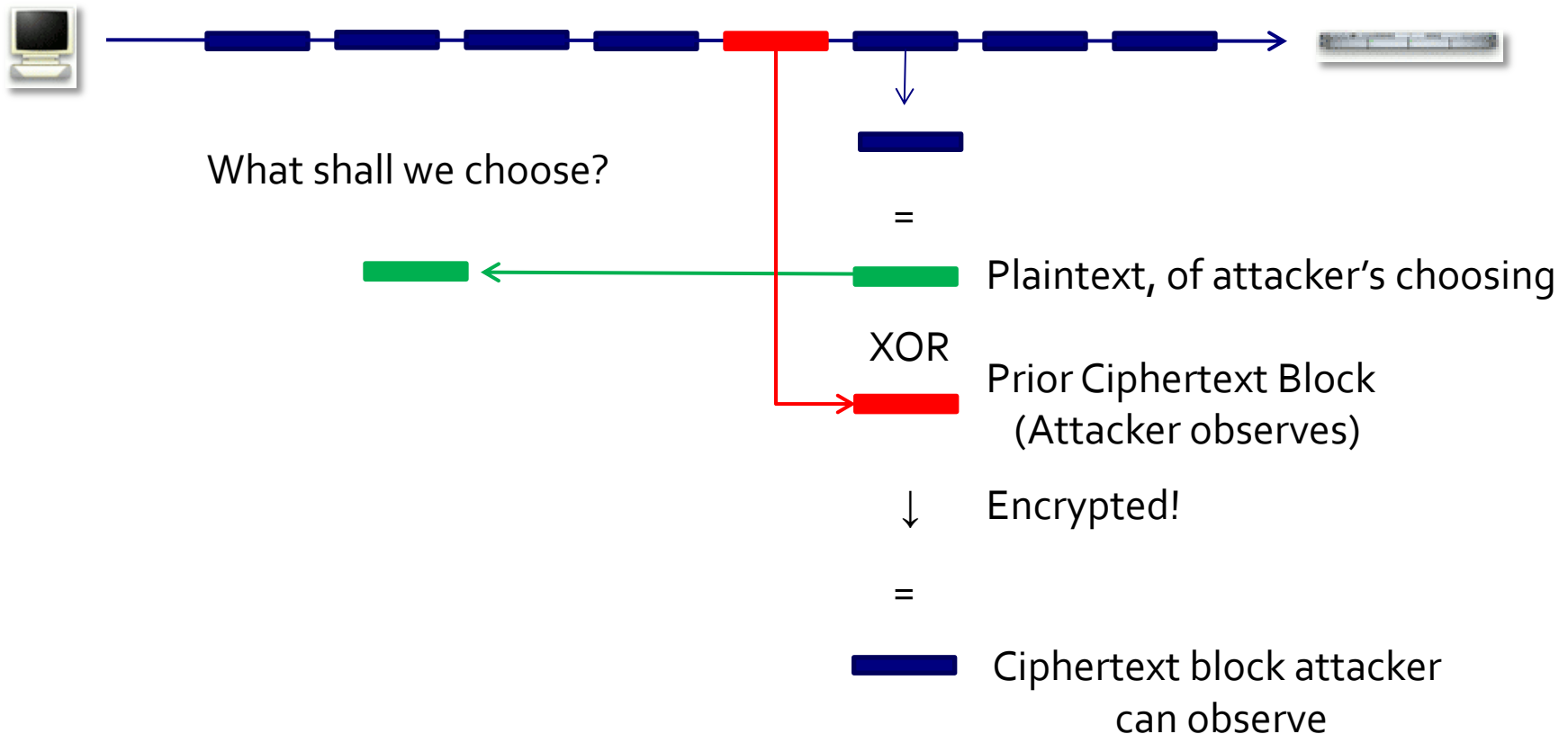
CBC Attack in SSL



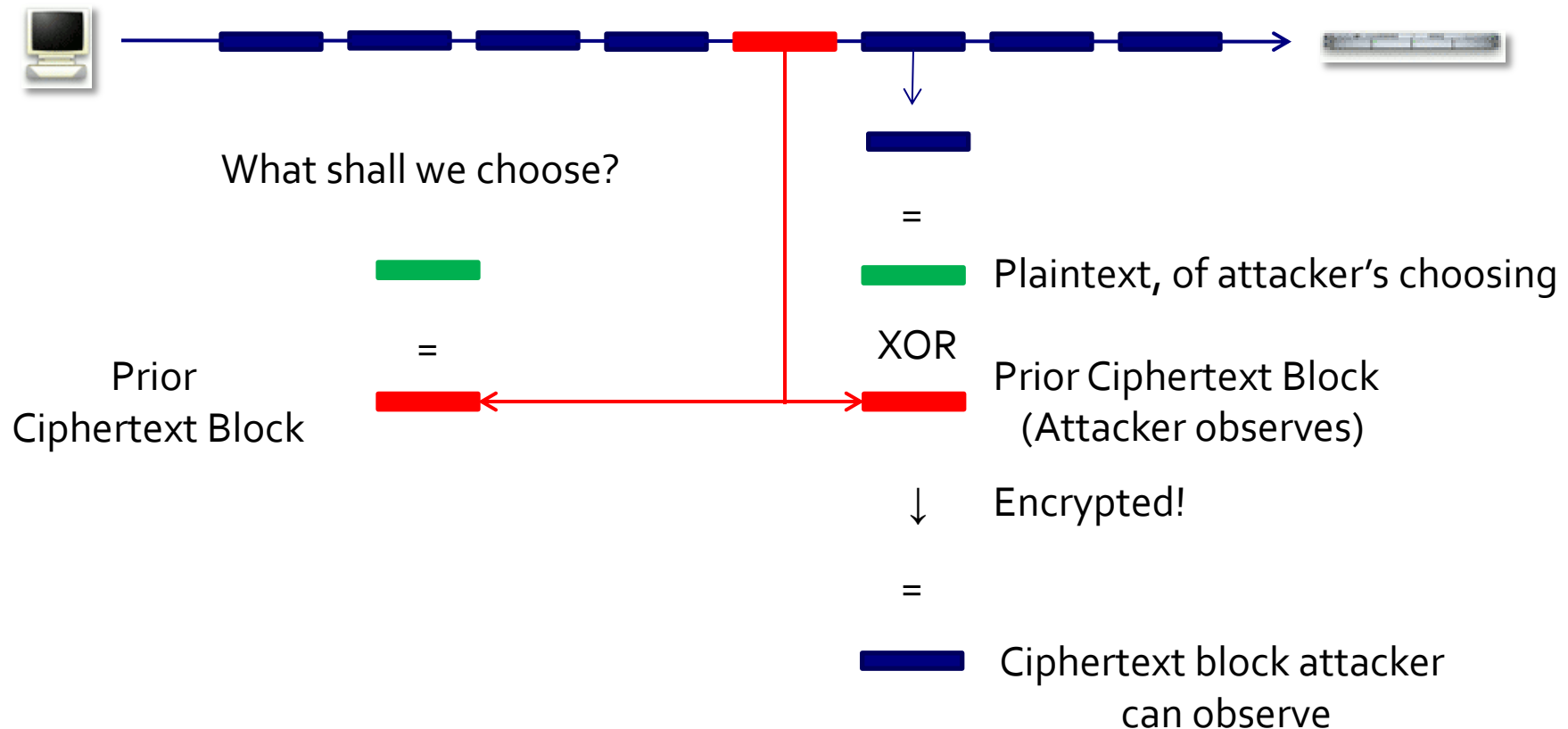
CBC Attack in SSL



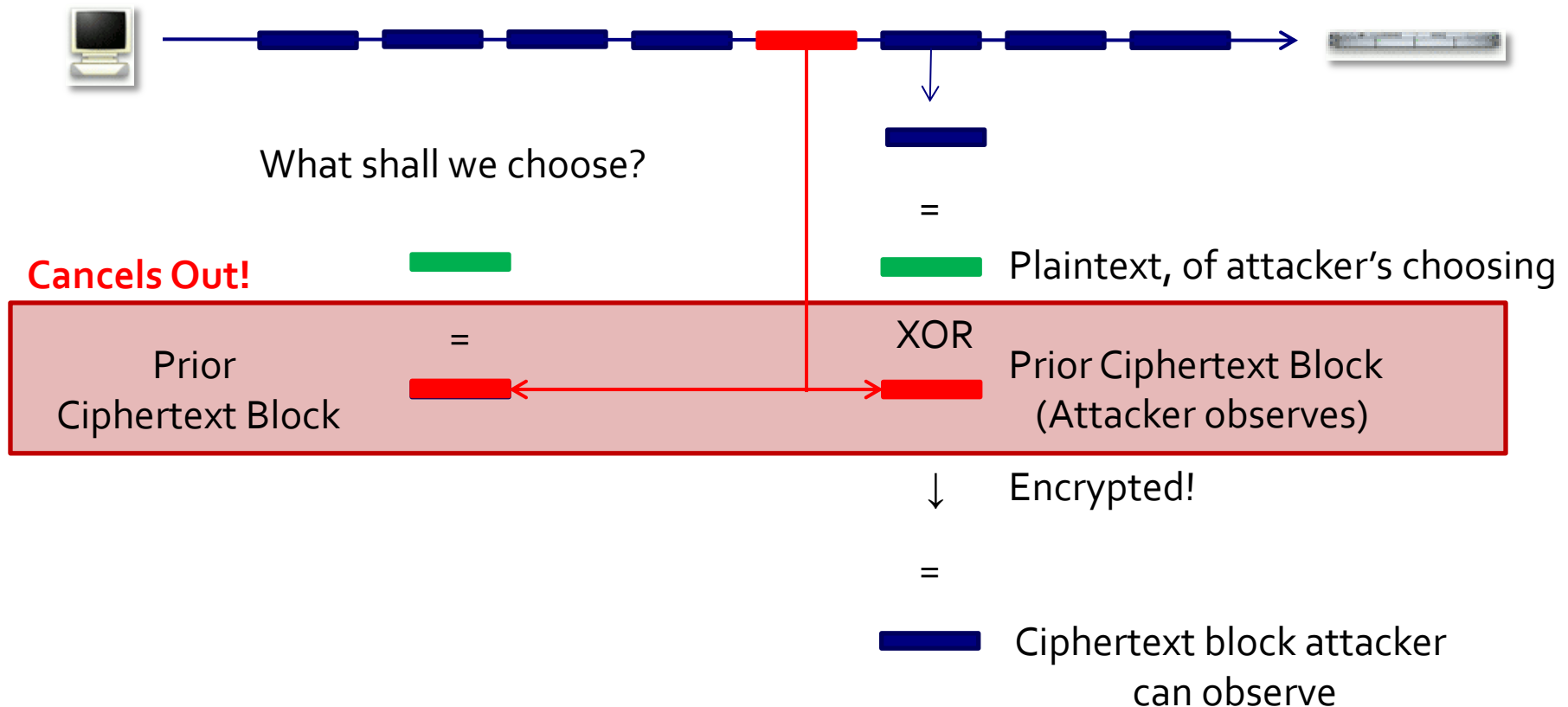
CBC Attack in SSL



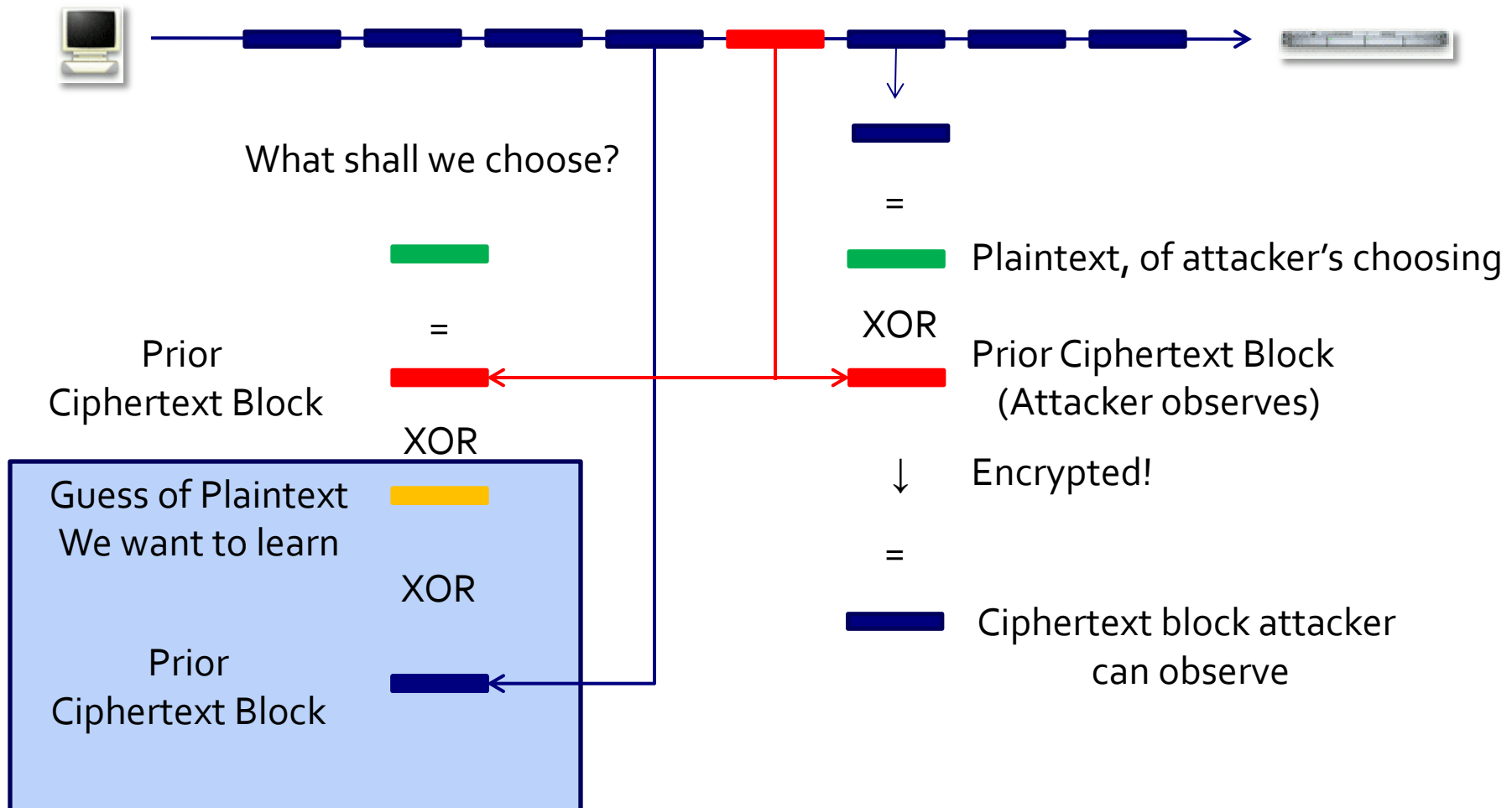
CBC Attack in SSL



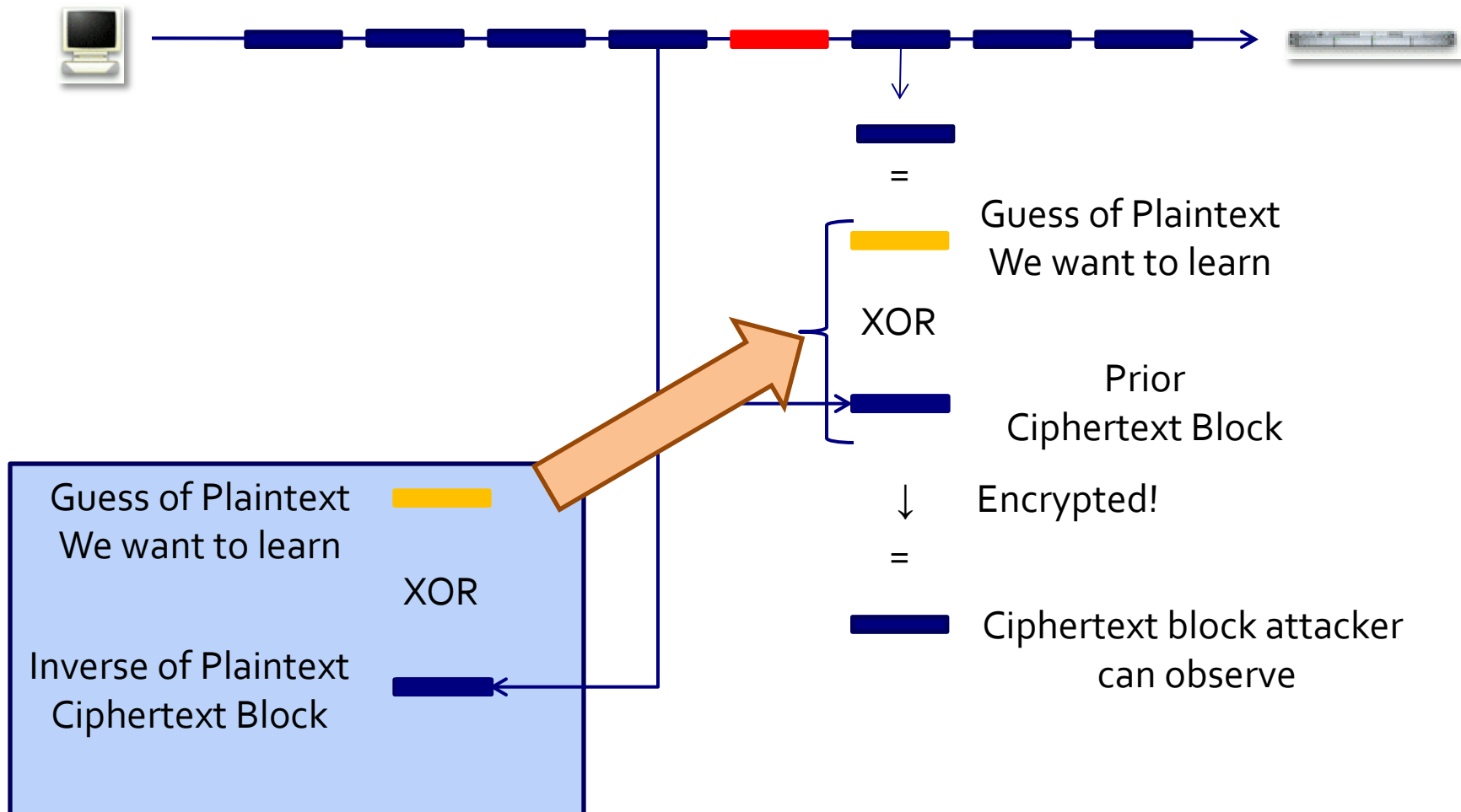
CBC Attack in SSL



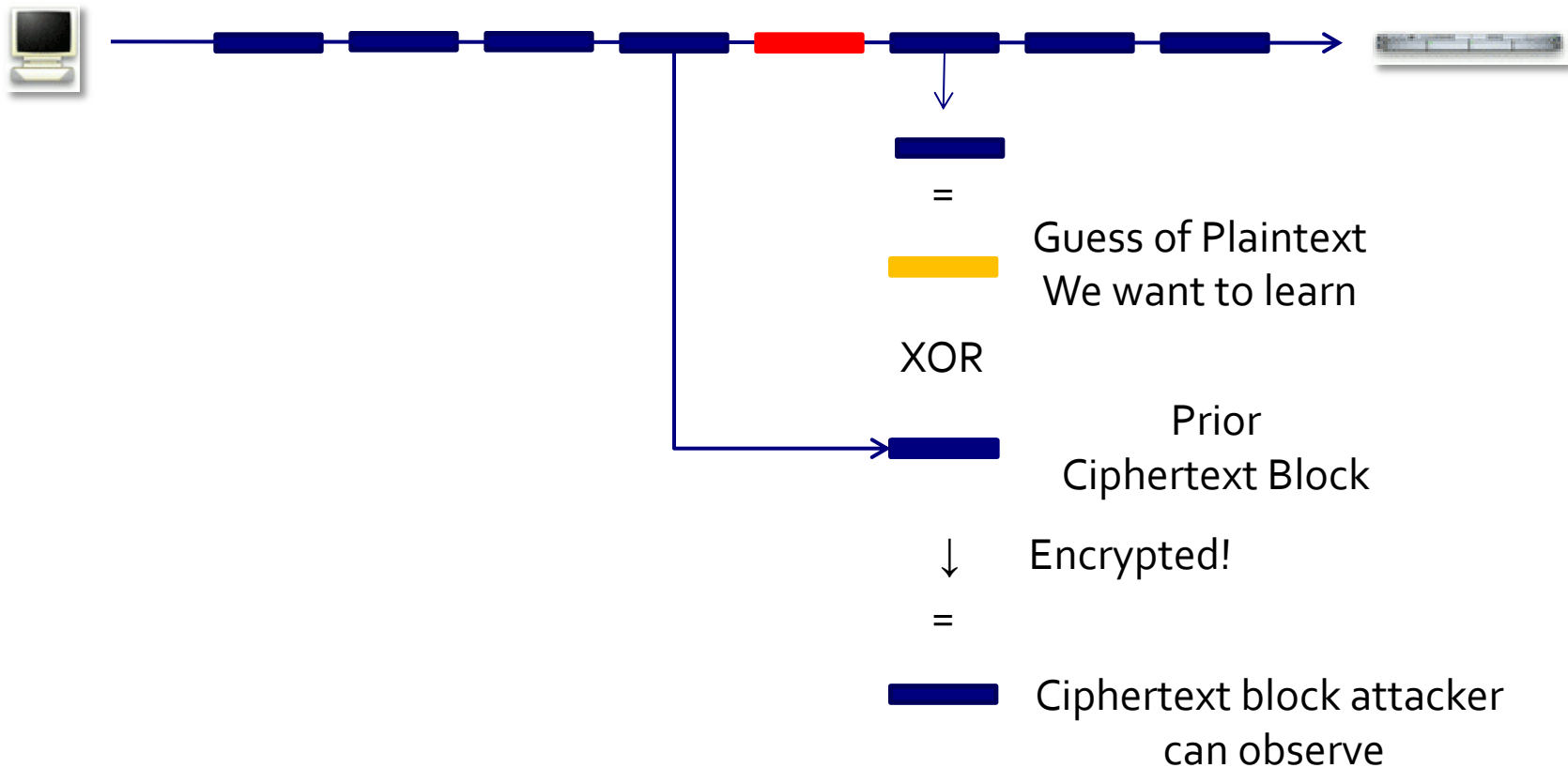
CBC Attack in SSL



CBC Attack in SSL

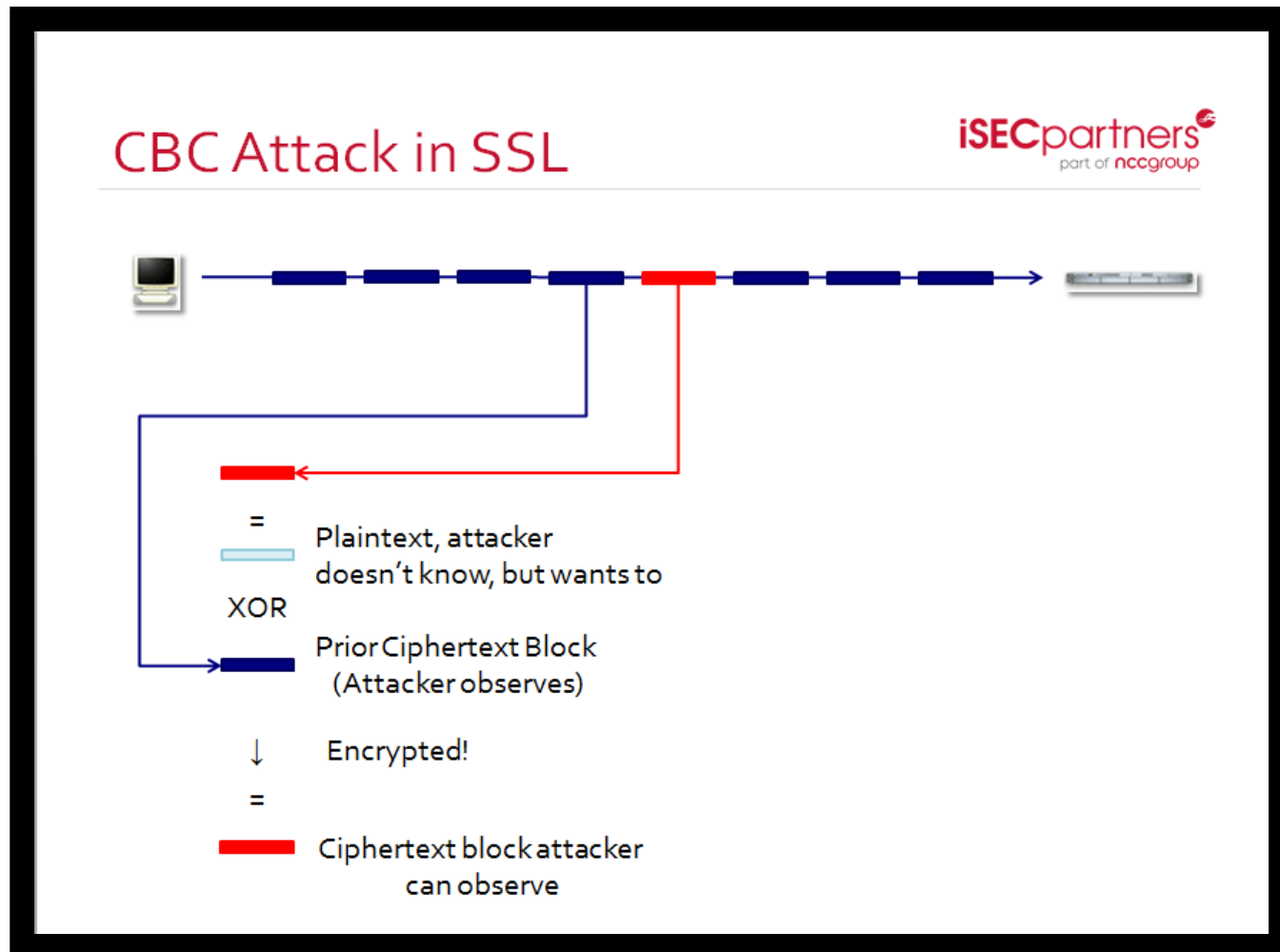


CBC Attack in SSL



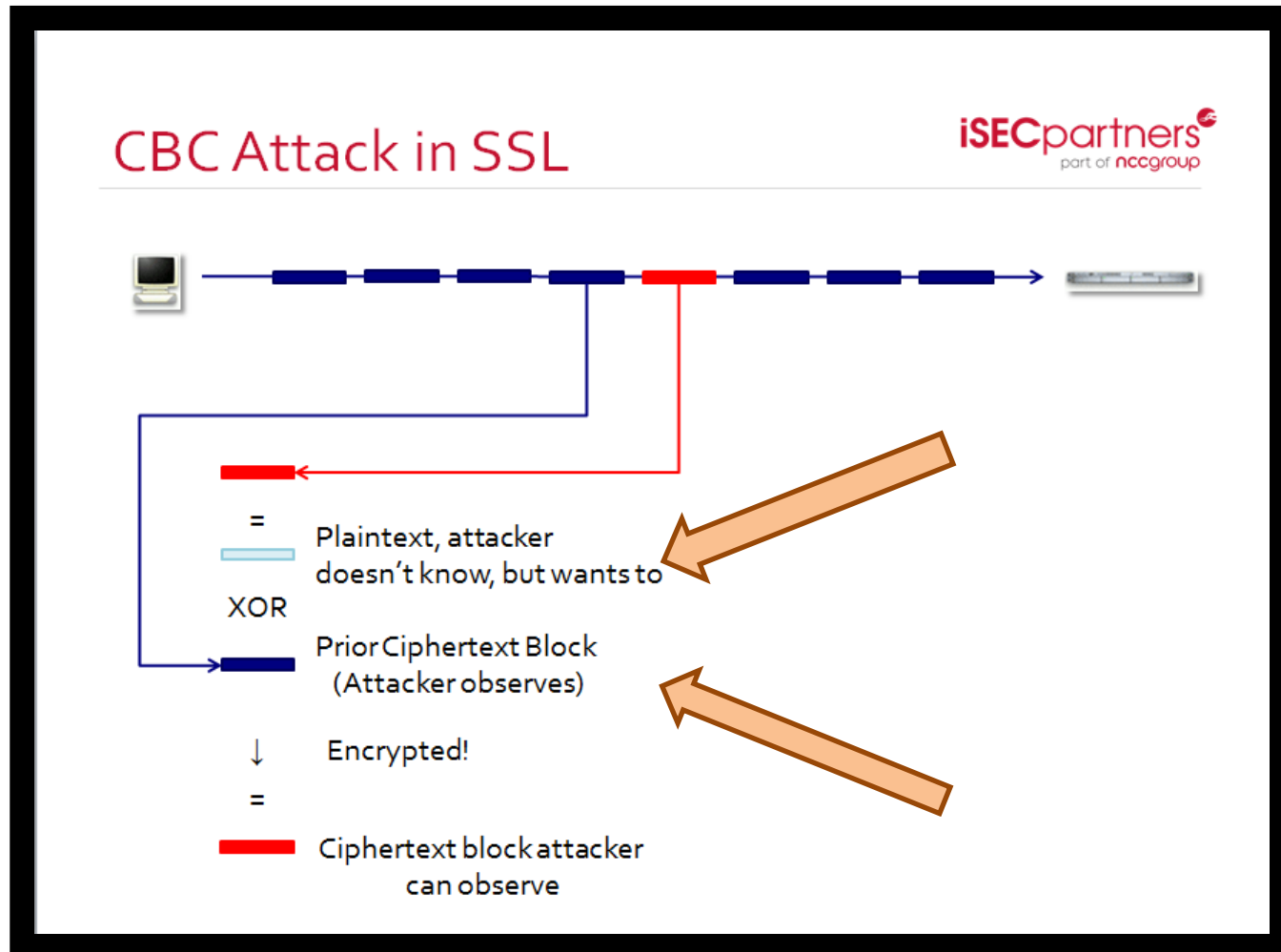
CBC Attack in SSL

- Remember This Slide?

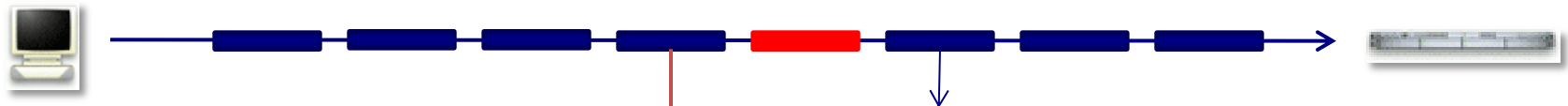


CBC Attack in SSL

- Remember This Slide?



CBC Attack in SSL



=

Guess of Plaintext
We want to learn



XOR



Prior
Ciphertext Block

↓

Encrypted!

=



Ciphertext block attacker
can observe

CBC Attack in SSL



=
Plaintext, attacker
doesn't know, but wants to

XOR

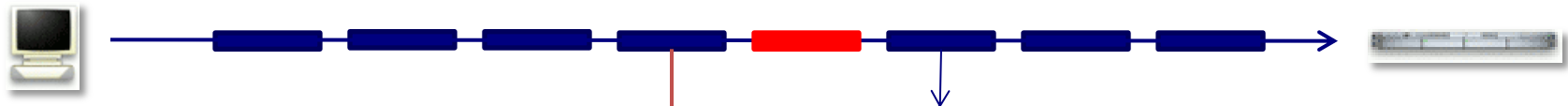
Prior Ciphertext Block
(Attacker observes)

↓
Encrypted!

=

Ciphertext block attacker
can observe

CBC Attack in SSL



=

Guess of Plaintext
We want to learn



XOR



Prior
Ciphertext Block

↓

Encrypted!

=



Ciphertext block attacker
can observe

CBC Attack in SSL



=
Plaintext, attacker
doesn't know, but wants to

XOR

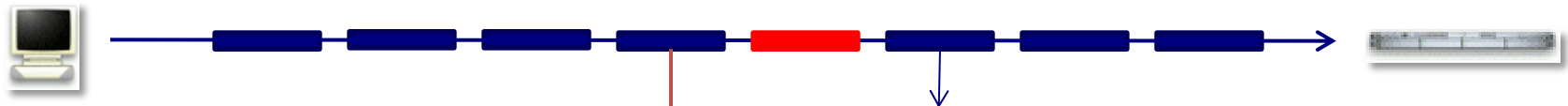
Prior Ciphertext Block
(Attacker observes)

↓
Encrypted!

=

Ciphertext block attacker
can observe

CBC Attack in SSL



=

Guess of Plaintext
We want to learn



XOR



Prior
Ciphertext Block

↓

Encrypted!

=



Ciphertext block attacker
can observe

CBC Attack in SSL



=
Plaintext, attacker
doesn't know, but wants to

XOR

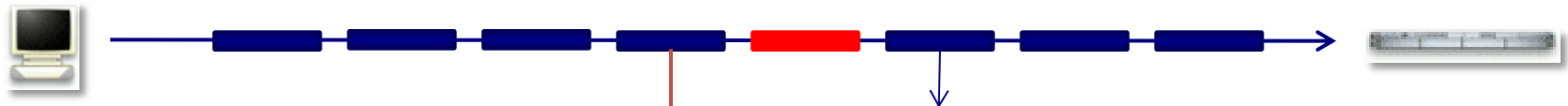
Prior Ciphertext Block
(Attacker observes)

↓
Encrypted!

=

Ciphertext block attacker
can observe

CBC Attack in SSL



CBC Attack in SSL



= Plaintext, attacker doesn't know, but wants to

XOR
Prior Ciphertext Block (Attacker observes)

↓ Encrypted!

=
Ciphertext block attacker can observe



=



XOR



Guess of Plaintext
We want to learn

Prior
Ciphertext Block

↓
Encrypted!

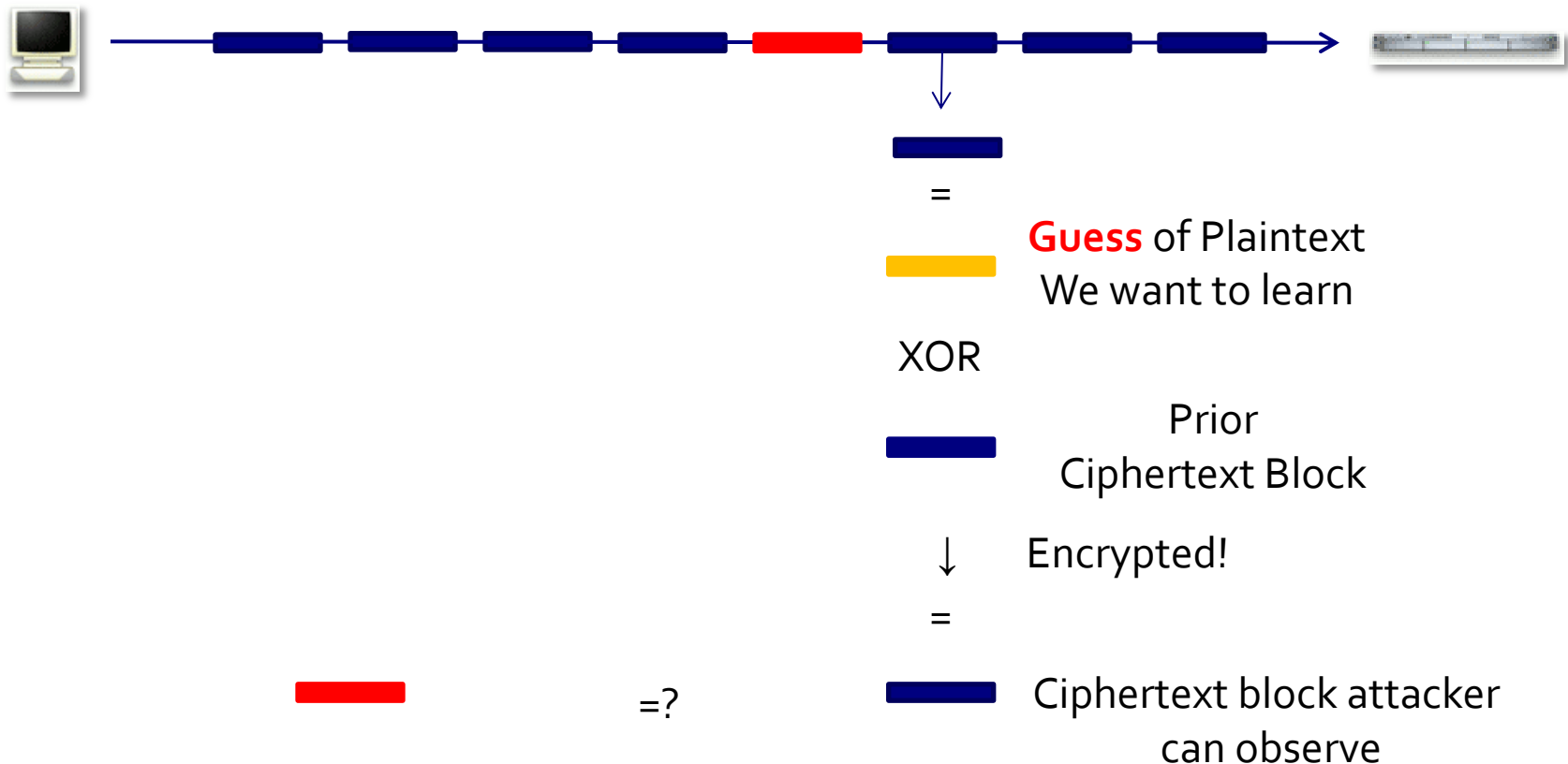
=



Ciphertext block attacker
can observe

If we guessed right, we will see an output that MATCHES a cipher text block we saw previously!

CBC Attack in SSL



If we guessed right, we will see an output that **MATCHES** a cipher text block we saw previously!

HTTP Request – What Do I Know?

POST /login HTTP/1.1

Host: bank.com

User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)
Gecko/20100101 Firefox/16.0

Cookie: a=secrets298fc1c149afbf4c8996fb924



I know everything but the cookie!

Chosen Boundary in a Slide

[illegible]

Chosen Boundary in a Slide

What the attacker knows

What the attacker DOESN'T know.

Different

Ciphertext

Blocks

Attacker pads to a block boundary!

G	E	T	/	...	C	O	k	i	e	:	a	=	s	e	c	r	e	t	...	s
---	---	---	---	-----	--------------	--------------	---	---	---	---	---	---	---	---	---	---	---	---	-----	---

A diagram of a 1D lattice with 24 sites. Sites 1-6 and 8-24 have downward arrows. Site 7 is a double site with an orange background and a diagonal line.

G	E	T		/	A	...	1	.	1	\	\	C	o	o	k	i	e	:		a	=	s		e
---	---	---	--	---	---	-----	---	---	---	---	---	---	---	---	---	---	---	---	--	---	---	---	--	---

A diagram consisting of 25 vertical lines. The first 6 lines have a downward arrow, followed by 6 empty lines, then 10 lines with downward arrows, and finally 3 empty lines and 1 line with a downward arrow.

Put It Together: BEAST

- Attacks SSL 3.0 TLS 1.0 with CBC Cipher suites
- Steals Cookies
- Works on HTTPS-only sites
 - Sorry Paypal

BEAST – Feasibility

- Ability to eavesdrop on the network
- Force victim to visit attackers page
- Ability to inject plaintext in an active SSL/TLS session

BEAST – Counter Measures

- Upgrade browsers
- Enable TLS 1.1, preferably 1.2
- Use RC4

Interlude: Protocol Downgrades

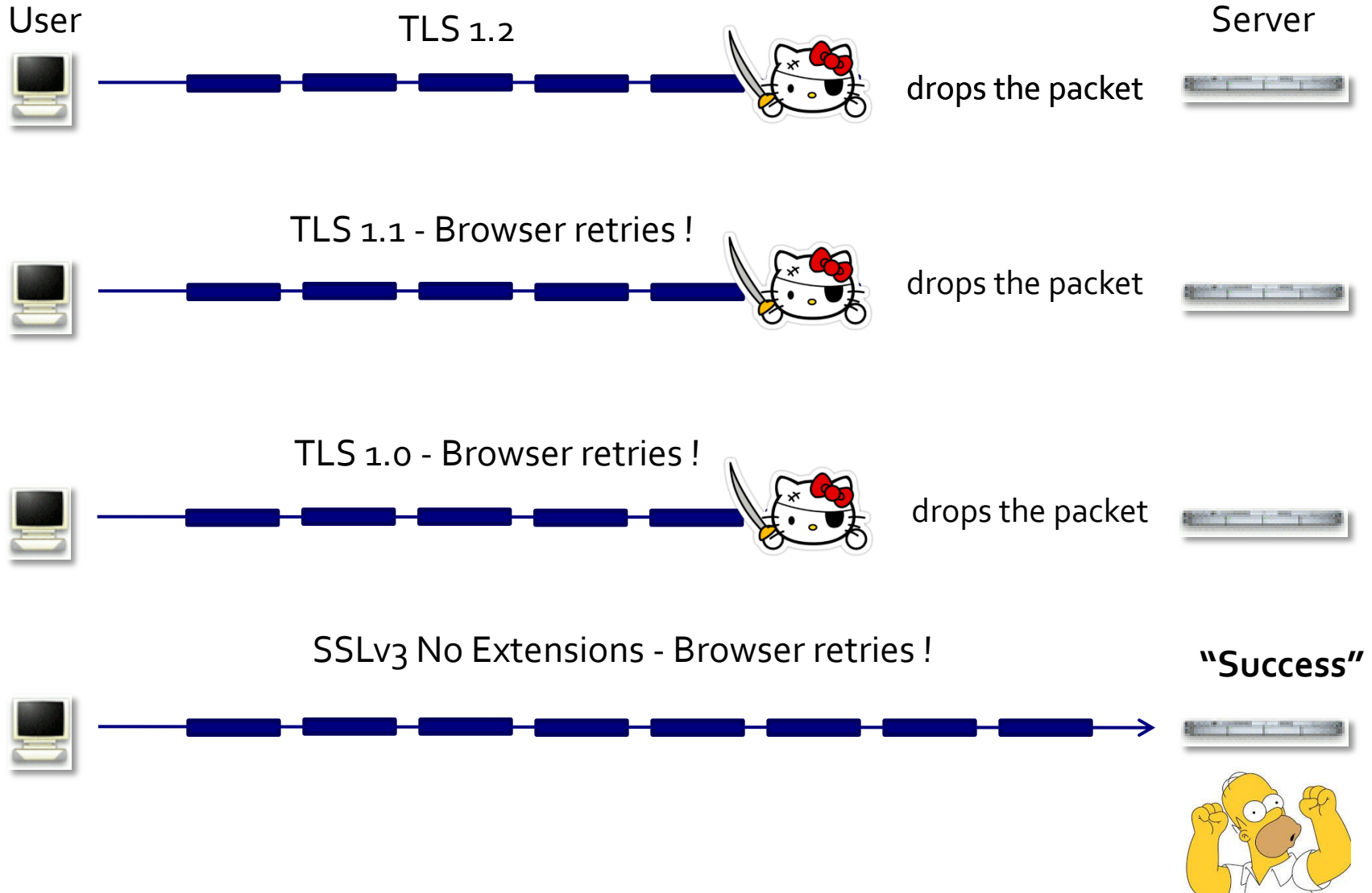
- We've mentioned TLS 1.1 and 1.2
- They're great!
- There's a problem:

Interlude: Protocol Downgrades

- We've mentioned TLS 1.1 and 1.2
- They're great!
- There's a problem:

They provide no security at all against an active attacker

Interlude: Protocol Downgrades



Interlude: Protocol Downgrades

Why Do Browsers Support Fallback?

- Networks Are Hostile to TLS 1.1+
 - Middleboxes don't recognize it and choke
- Sites Can't Speak TLS 1.1+
 - Sometimes an error (not so bad)
 - Sometimes they just hang (quite bad)

Interlude: Protocol Downgrades

- Until Browsers Remove Fallback to TLS 1/SSLv3 we cannot fully rely on TLS 1.1+
- Until sites stop breaking for TLS 1.1+ Browsers can't Remove the Fallback
- Not to call anyone out.... But....
 - <https://www.imperialviolet.org/2013/10/07/f5update.html>

Lucky 13

**DO YOU
FEEL
LUCKY....
PUNK?**

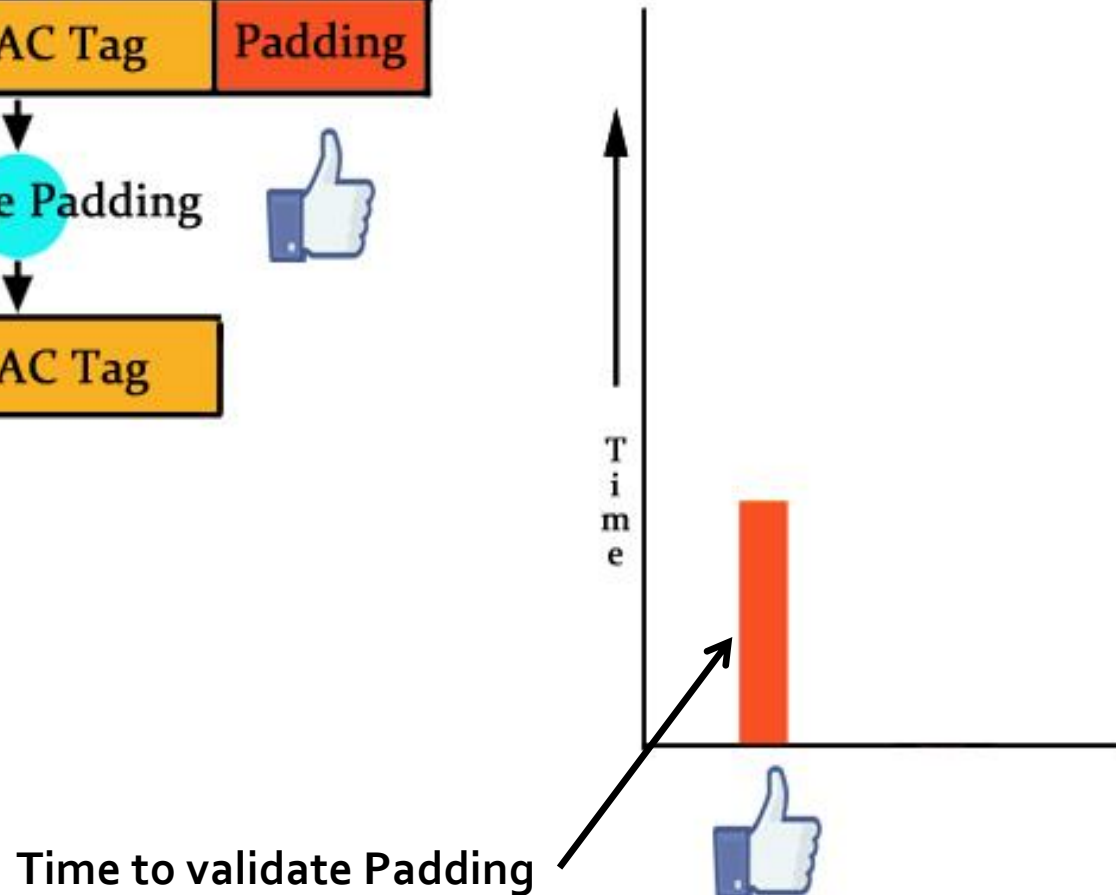
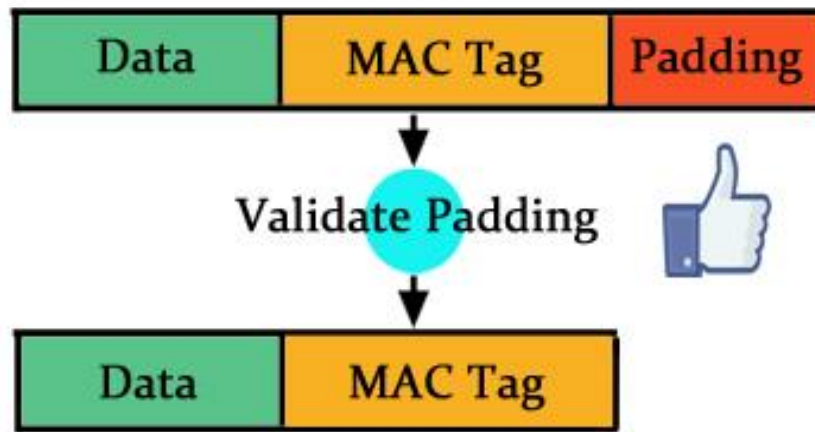


**WELL....
DO YA?**

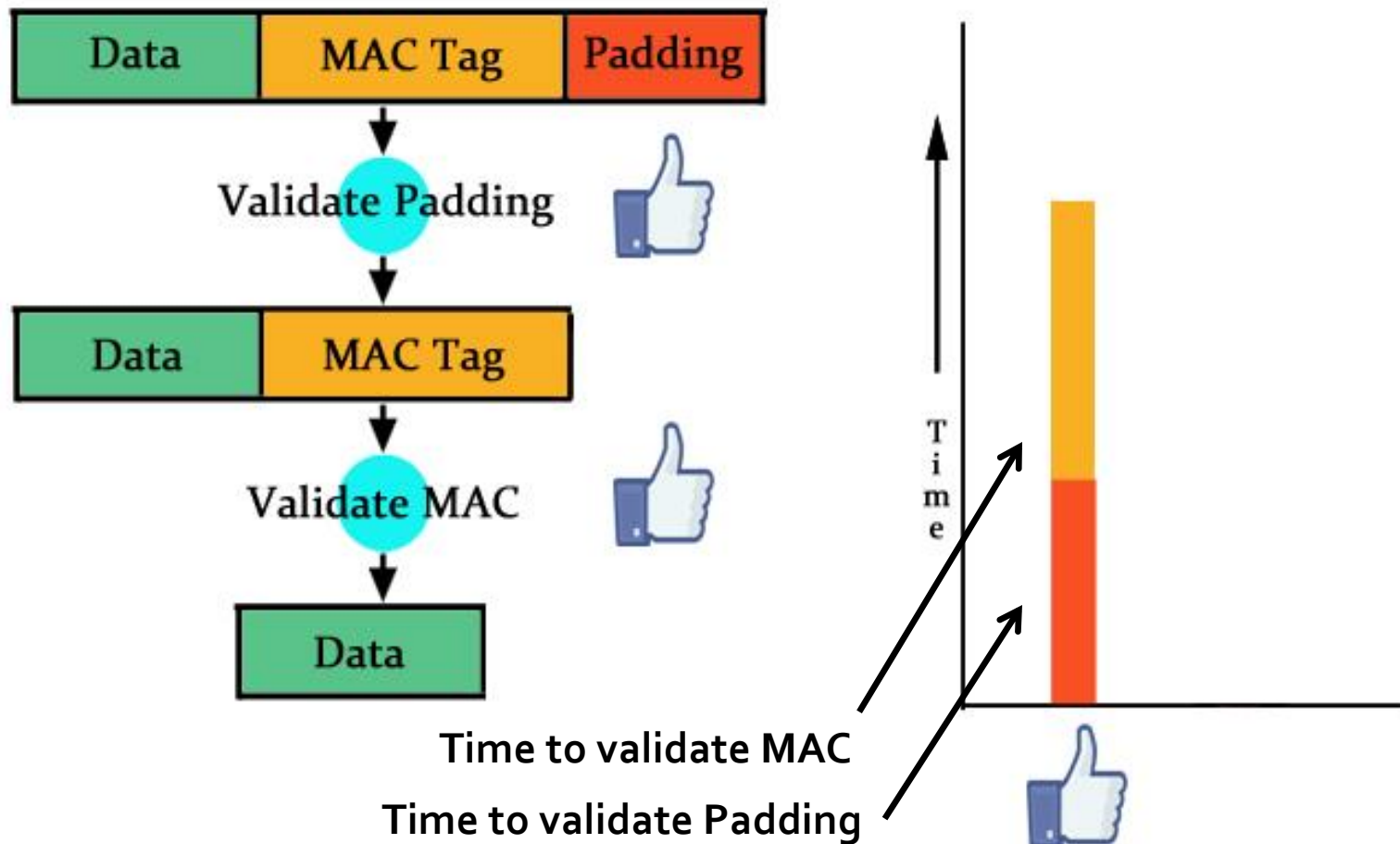
Lucky 13 - What is it ?

- Successor of Padding Oracle Attack
- Timing attack on CBC encryption mode
- 13 bytes of header information in TLS MAC calculation leaks timing information during decryption.

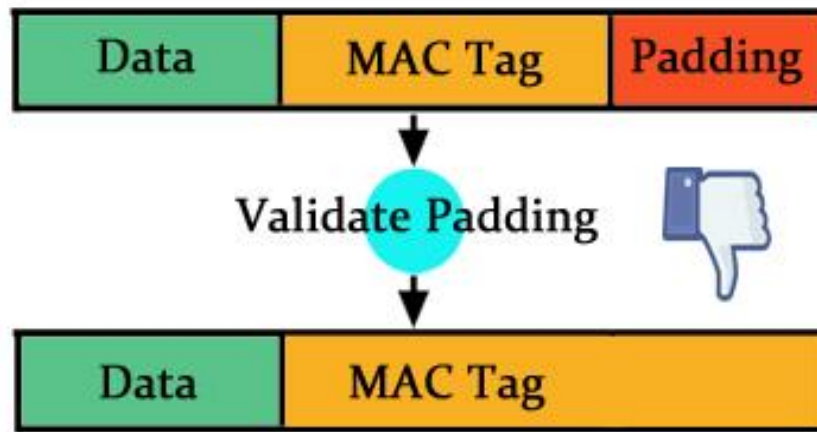
Lucky 13 – How it works



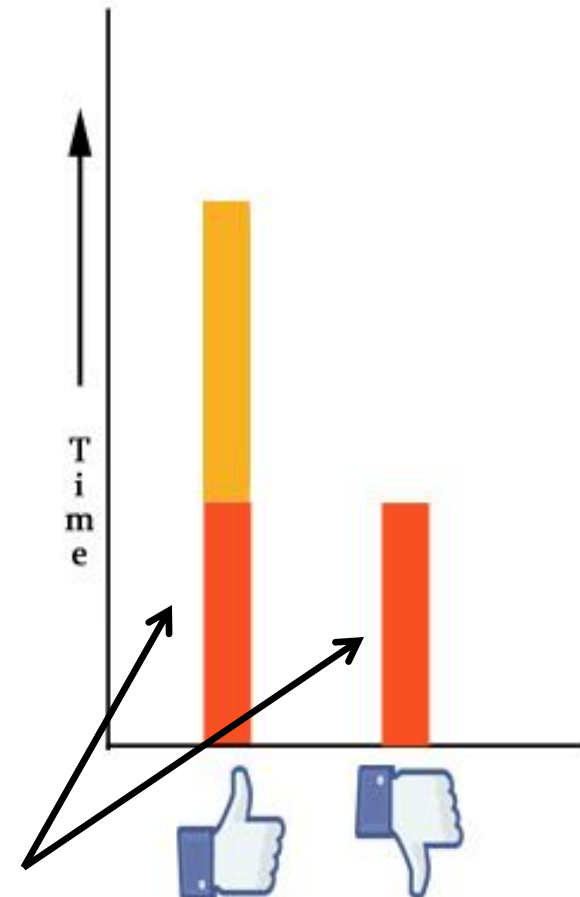
Lucky 13 – How it works



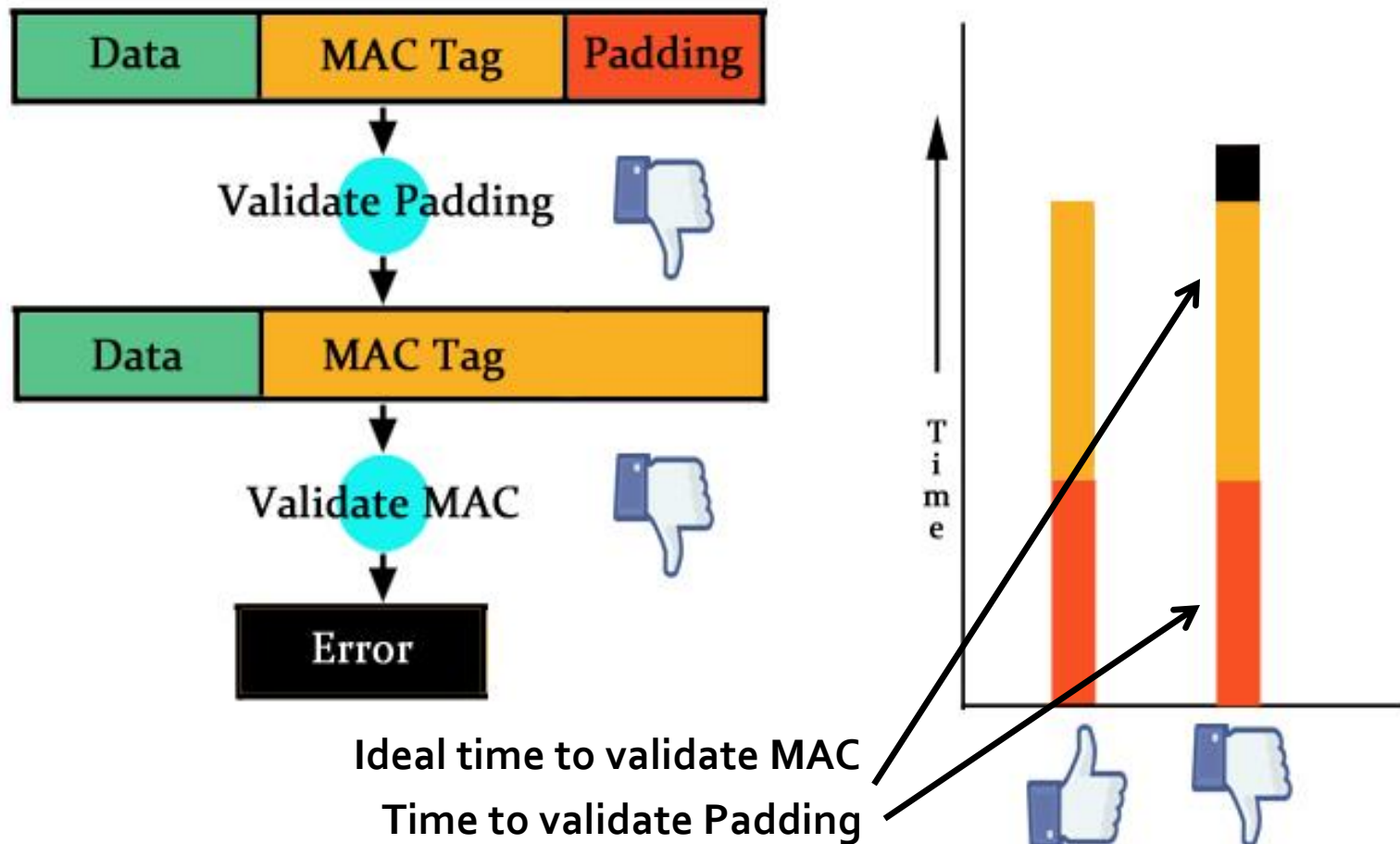
Lucky 13 – How it works



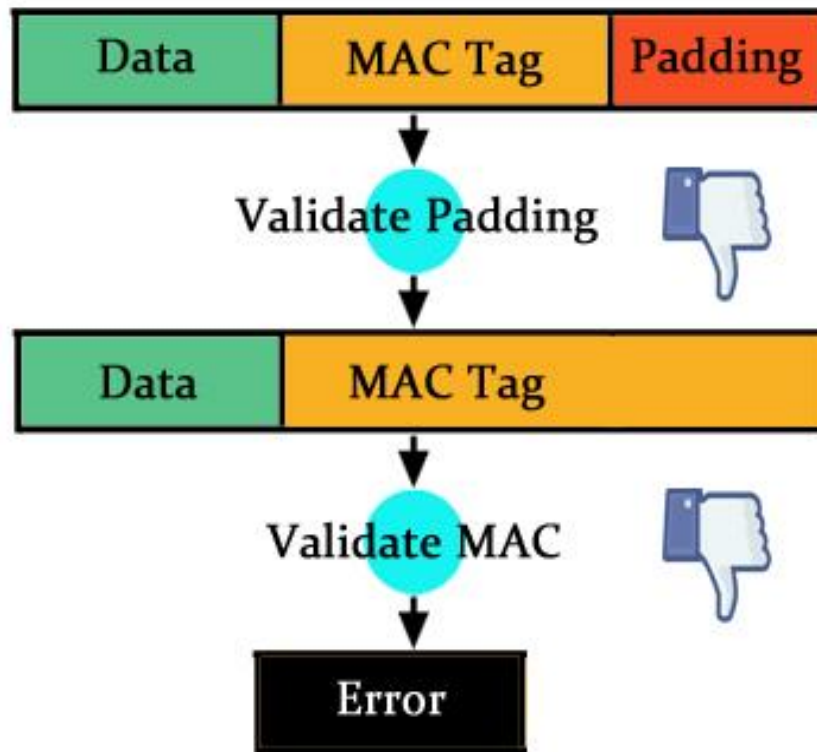
Time to validate Padding



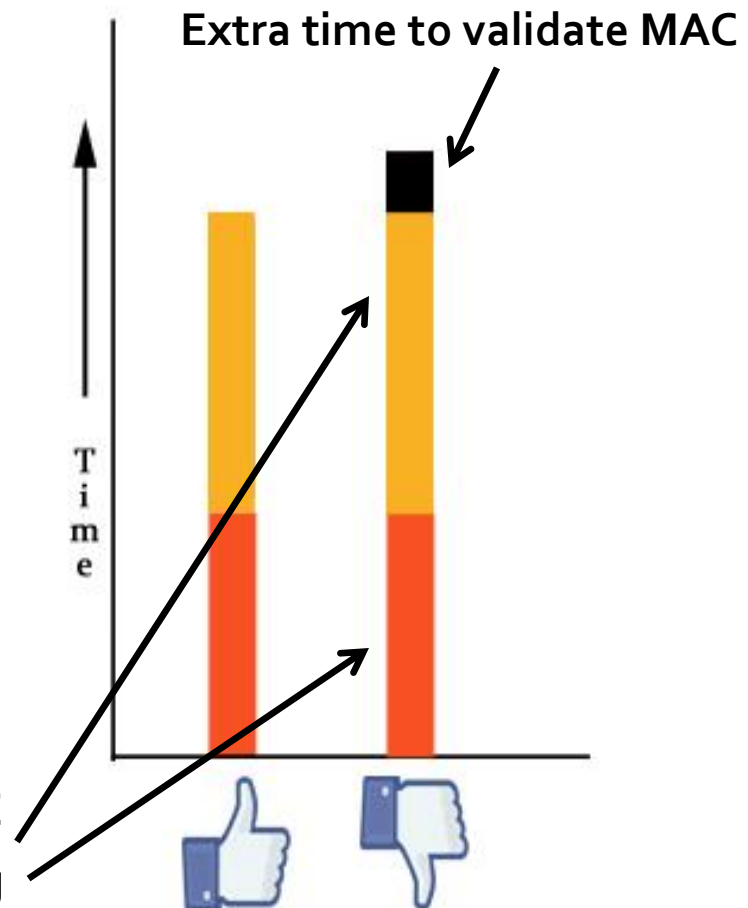
Lucky 13 – How it works



Lucky 13 – How it works



Ideal time to validate MAC
Time to validate Padding



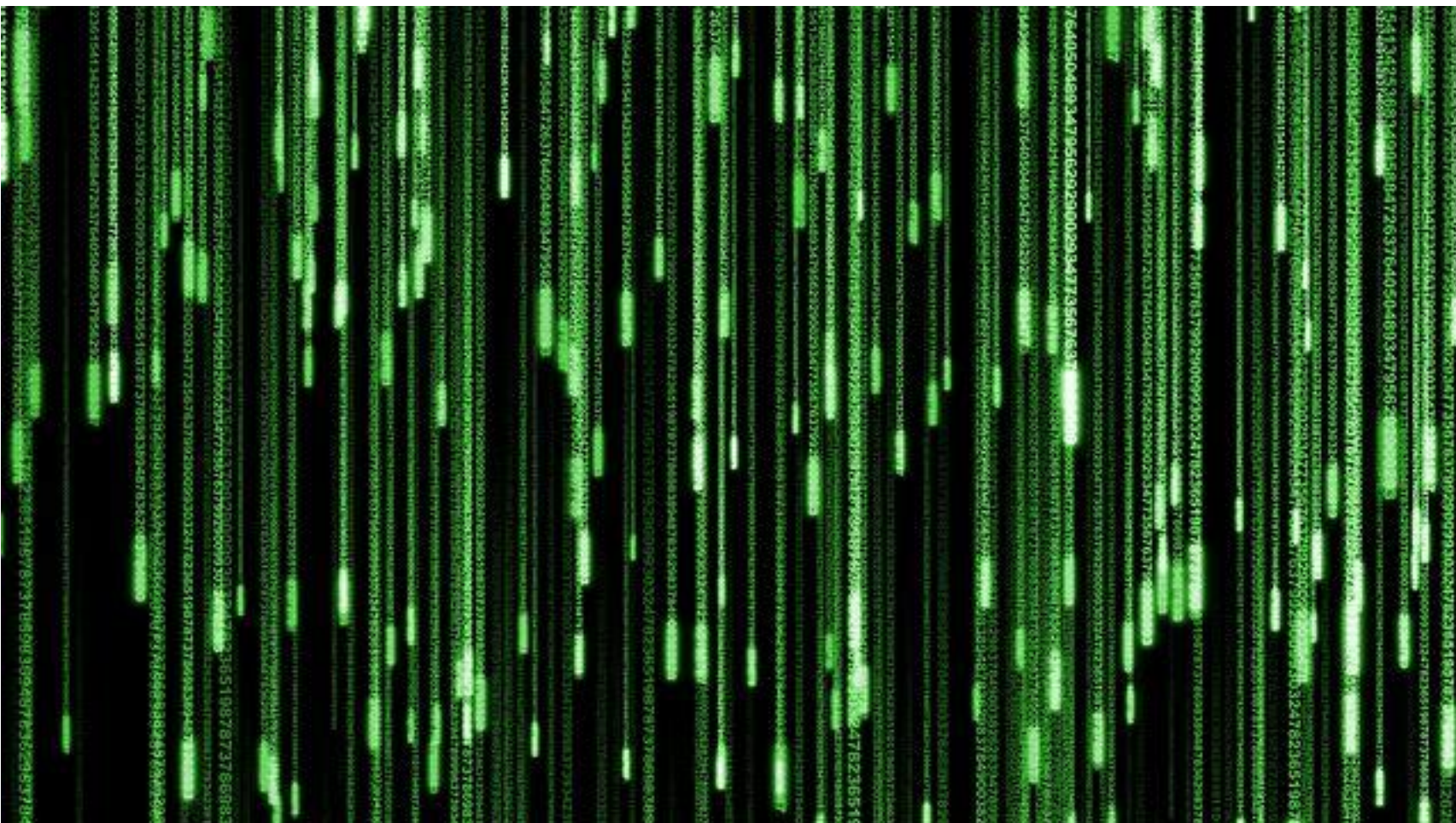
Lucky 13 – Feasibility

- Needs Man-in-the-middle
- CBC-mode encryption in versions of TLS are potentially vulnerable.
- Requires huge number of request
- Requires no Network jitter

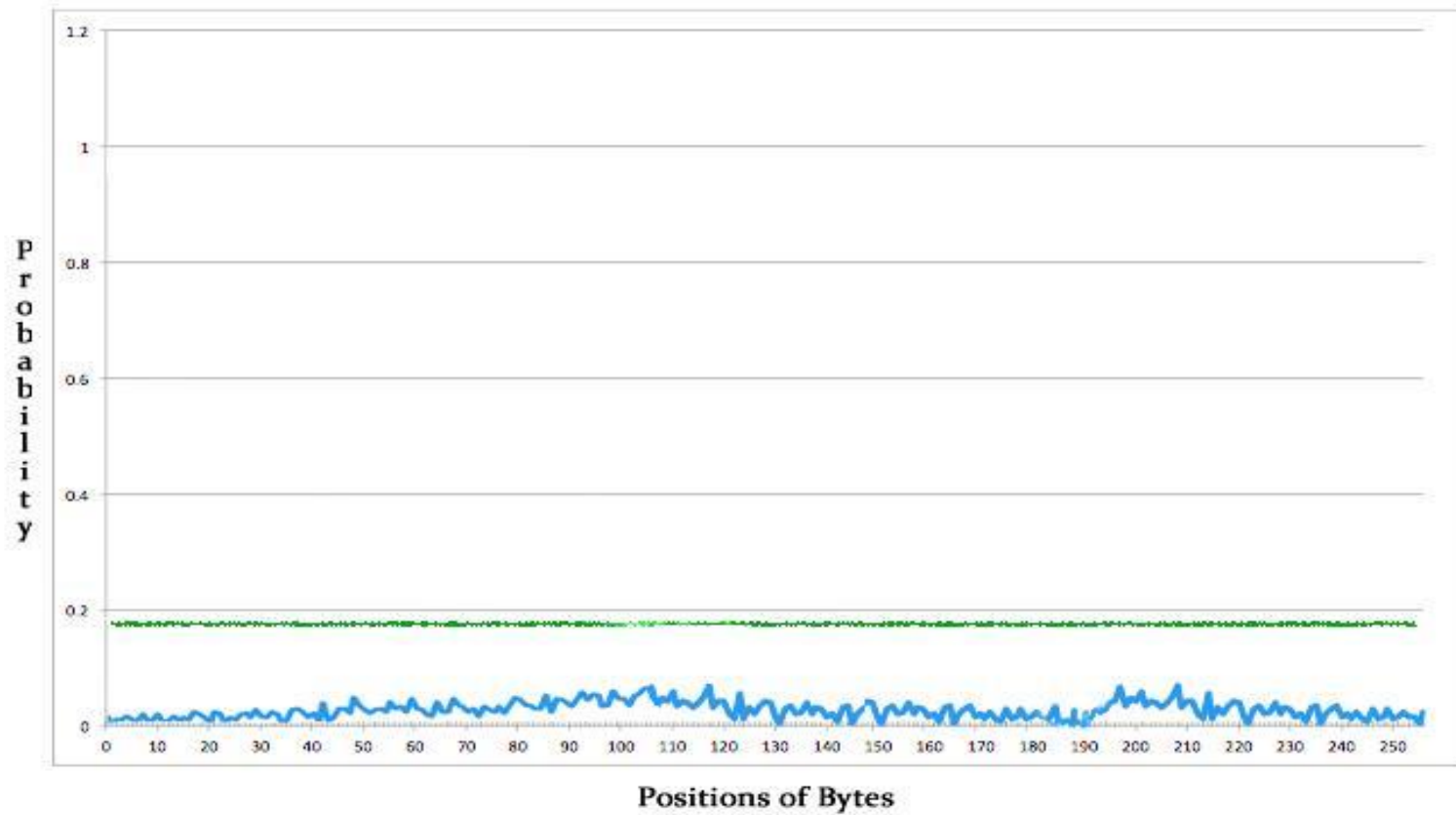
Lucky 13 – Counter Measures

- Uniform processing time to decrypt ciphertexts
- Add random timing delays to the decryption for any timing attack
- Using stream cipher like RC4
- Using an authenticated encryption algorithm, such as AES-GCM

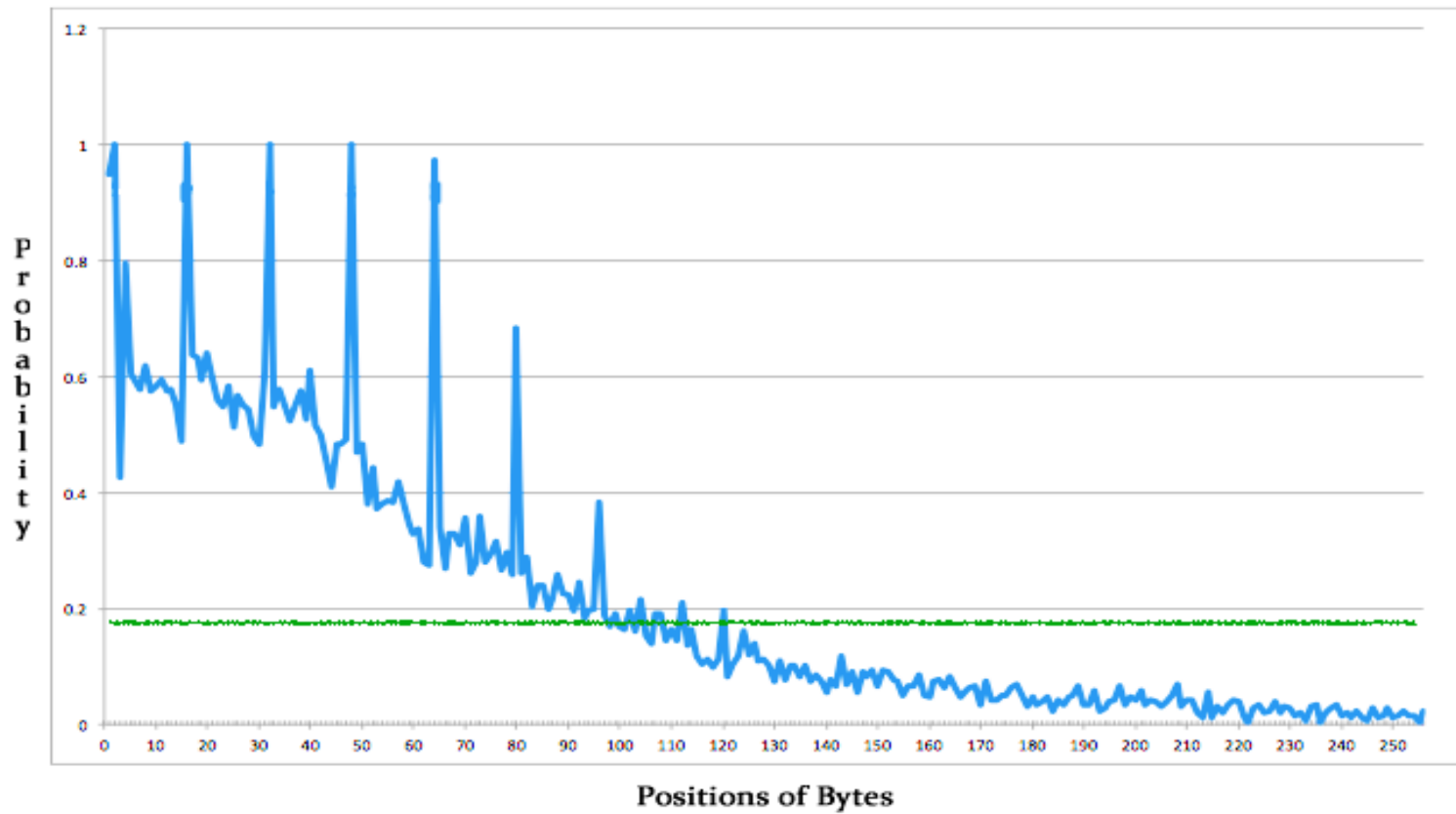
RC4



RC₄ Biases - What is it ?



RC₄ Biases - What is it ?



RC₄ Biases – Feasibility

- Force victim to renegotiate.
- This attack will require over **4 billion** SSL connections or re-negotiations for an individual HTTP session.

RC4 Biases – Counter Measures

- Researchers still working on finding mitigations of this issue.
- Temporary mitigations
 - Throttle client initiated re-negotiations and connections from individual IP addresses
 - If possible use block ciphers with mitigations of timing and CBC mode encryption attack mitigated

Compression



How compression works?

- DEFLATE compression mechanism

Etiam euismod. Fusce facilisis lacinia dui. Suspendisse potenti. In mi erat, cursus id, nonummy sed, ullamcorper eget, sapien. Praesent pretium, magna in eleifend egestas, pede pede pretium lorem, quis consectetur tortor sapien facilisis magna.

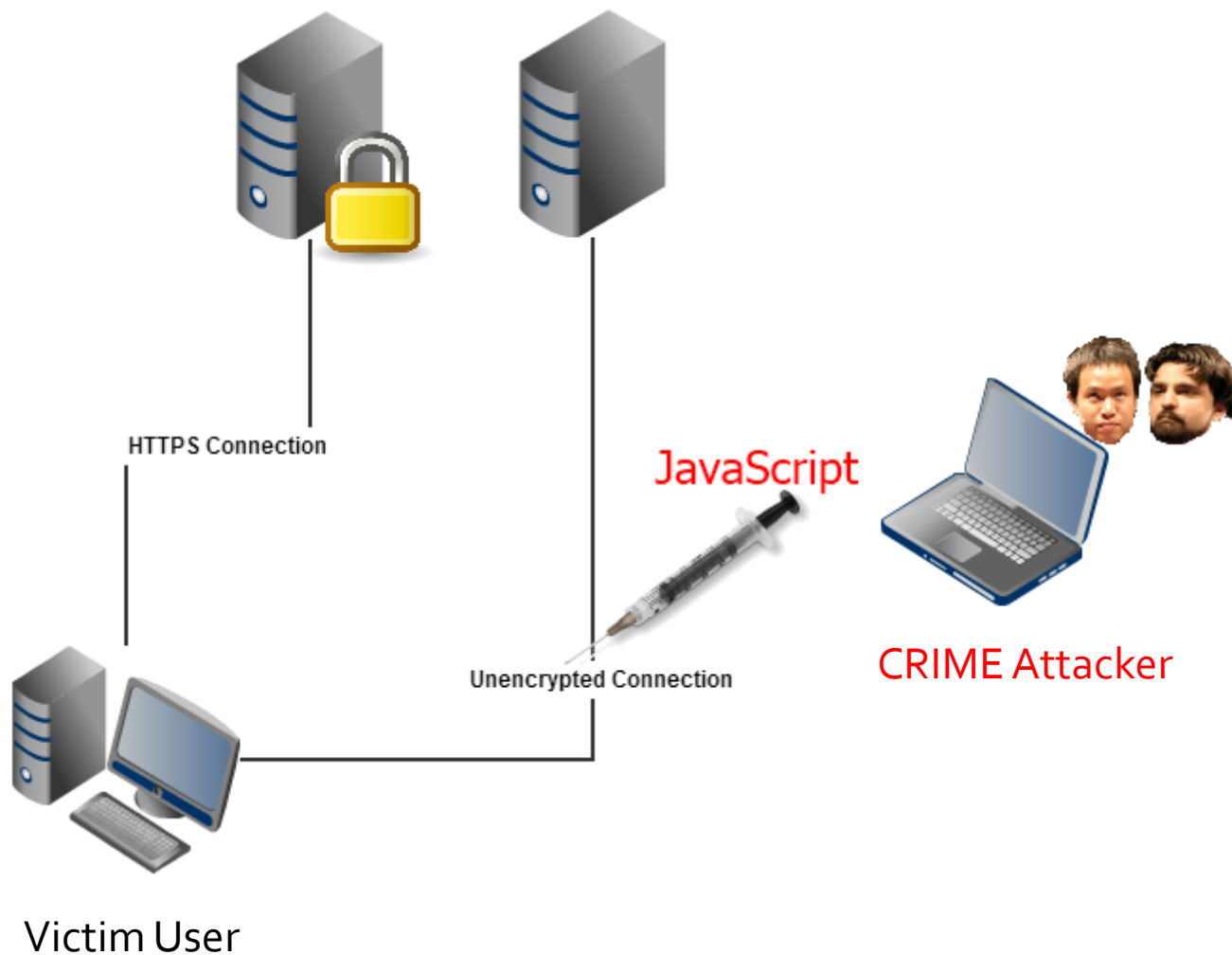
CRIME



CRIME - What is it ?

- **C**ompression **R**atio **I**nfo-leak **M**ade **E**asy
- Chosen plaintext attack on HTTP request
- Uses size information in TLS compression to recover plaintext cookies

CRIME – How it works



CRIME – How it works

```
GET /evil_request_path HTTP/1.1
Host: bank.com
User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)
Gecko/20100101 Firefox/16.0
Cookie: sessionid=d3b0c44298fc1c149afb4c8996fb924
```



**Attacker doesn't control entire request,
but can see its cipher text on the wire**

CRIME – How it works

GET /evil_request_path HTTP/1.1

Host: bank.com

User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)

Gecko/20100101 Firefox/16.0

Cookie: sessionid=d3b0c44298fc1c149afbf4c8996fb924



Attacker fully controls request path

CRIME – How it works

GET /evil_request_path HTTP/1.1

Host: bank.com

User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)

Gecko/20100101 Firefox/16.0

Cookie: sessionId=d3b0c44298fc1c149afb4c8996fb924



Attacker does not see, but can infer these values

CRIME – How it works

GET /evil_request_path HTTP/1.1

Host: bank.com

User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)

Gecko/20100101 Firefox/16.0

Cookie: sessionid=d3b0c44298fc1c149afb4c8996fb924



Attacker cannot see/control, wants to steal

CRIME – How it works

GET /**sessionid=a** HTTP/1.1

Host: bank.com

User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)
Gecko/20100101 Firefox/16.0

Cookie: **sessionid=d3b0c44298fc1c149afb4c8996fb924**

=> Compressed Length = 12,494 bytes – **Not a match**

GET /**sessionid=d** HTTP/1.1

Host: bank.com

User-Agent: Mozilla/5.0 (Windows NT 6.1; WOW64; rv:16.0)
Gecko/20100101 Firefox/16.0

Cookie: **sessionid=d3b0c44298fc1c149afb4c8996fb924**

=> Compressed Length = 12,493 bytes – **Possible match**

- The attacker can intercept the victim's network traffic.
- Victim authenticates to a website over HTTPS and negotiates TLS Compression with the server.
- Victim accesses a non-HTTPS website.
- Browser supporting TLS Compression

CRIME – Counter Measures

- Disabling TLS compression on both Browser and Server side.
- Updated Browser versions:
 - Chrome: 21.0.1180.89 and above
 - Firefox: 15.0.1 and above
 - Opera: 12.01 and above
 - Safari: 5.1.7 and above
- Apache 2.2 using mod_SSL:
SSLCompression flag is set to "*SSLCompression off*"
- Apache using mod_gnutls :
GnuTLSPriorities flag = "*!COMP-DEFLATE*"

BREACH



BREACH – What is it ?

- **B**rowser **R**econnaissance and **E**xfiltration via **A**daptive **C**ompression of **H**ypertext
- Chosen plaintext attack on HTTP response
- Uses difference of response size information in due to varying sizes of HTTP compression to recover plaintext secret information
- Resurrection of CRIME

BREACH – How it works

```
<form target="https://example.com:443/  
products/catalogue.aspx?id=12345&user=username" >  
...  
<td nowrap id="tdErrLgf">  
<a href="logoff.aspx?CSRFtoken=4bd634cda846fd7cb4cb00  
31ba249ca2">Log Off</a></td>
```

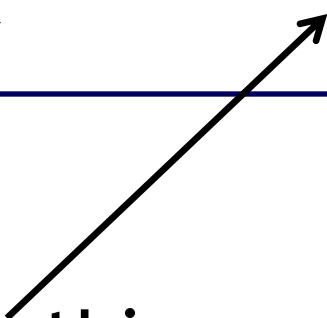


Attacker can control this value

BREACH – How it works

```
<form target="https://example.com:443/  
products/catalogue.aspx?id=12345&user=username" >  
...  
<td nowrap id="tdErrLgf">  
<a href="logout.aspx?CSRFtoken=4bd634cda846fd7cb4cb00  
31ba249ca2">Log Off</a></td>
```

Attacker cannot control this parameter, want to steal it



BREACH – How it works

```
GET /product/?id=12345&user=CSRFtoken=a HTTP /1.1  
Host: example.com
```

```
<form target="https://example.com:443/  
products/catalogue.aspx?id=12345&user=CSRFtoken=a" >  
...  
<td nowrap id="tdErrLgf">  
<a href="logoff.aspx?CSRFtoken=4bd634cda846fd7cb4cb00  
31ba249ca2">Log Off</a></td>
```

Size of response < Previous size = Match

Size of response >= Previous size = Mismatch

BREACH – Feasibility

- The application supports **HTTP compression**.
- The response should **reflect** back user's input.
- The response should have some **sensitive/ secret** information embedded in the body.

BREACH – Counter Measures

- Mask the secret:
 - new secret = random || (random \oplus previous secret)
- Enable anti-automation techniques
- Monitor your traffic
- Separate secrets from user input
- ~~• Disable HTTP compression~~

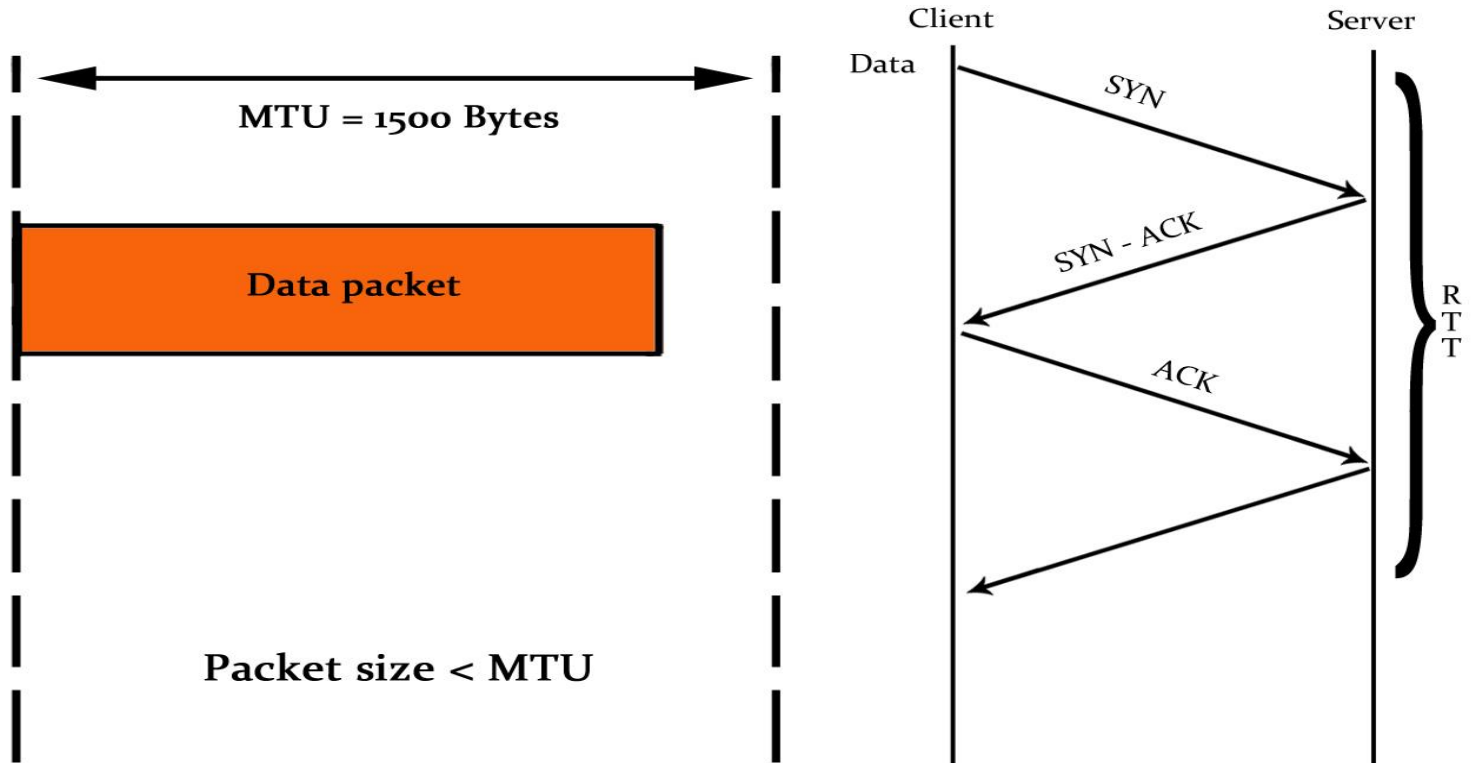
TIME



TIME - What is it ?

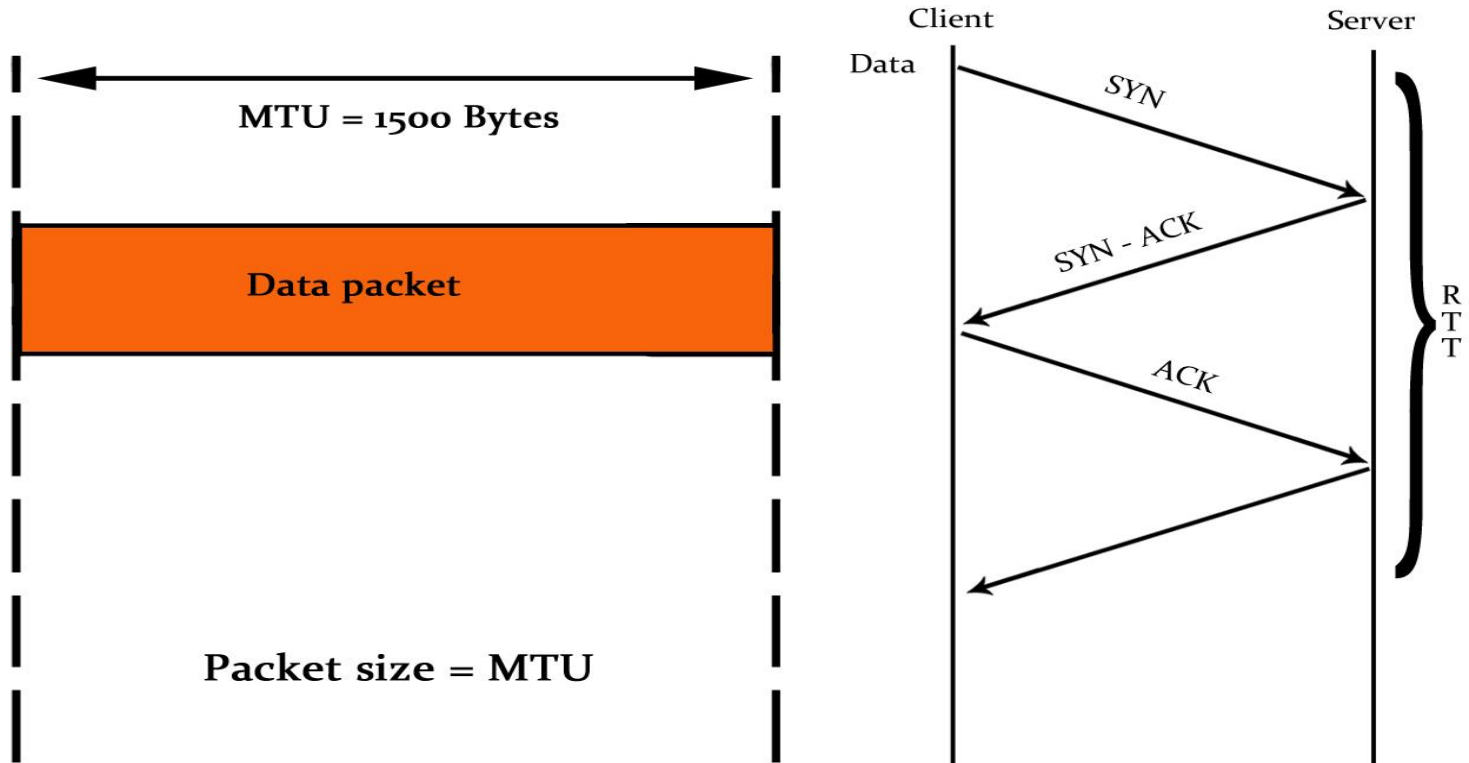
- Timing Info-leak **M**ade **E**asy
- Chosen plaintext attack on HTTP response
- Uses difference of response time information in due to varying sizes of HTTP compression to recover plaintext secret information
- Resurrection of CRIME

TIME – How it works



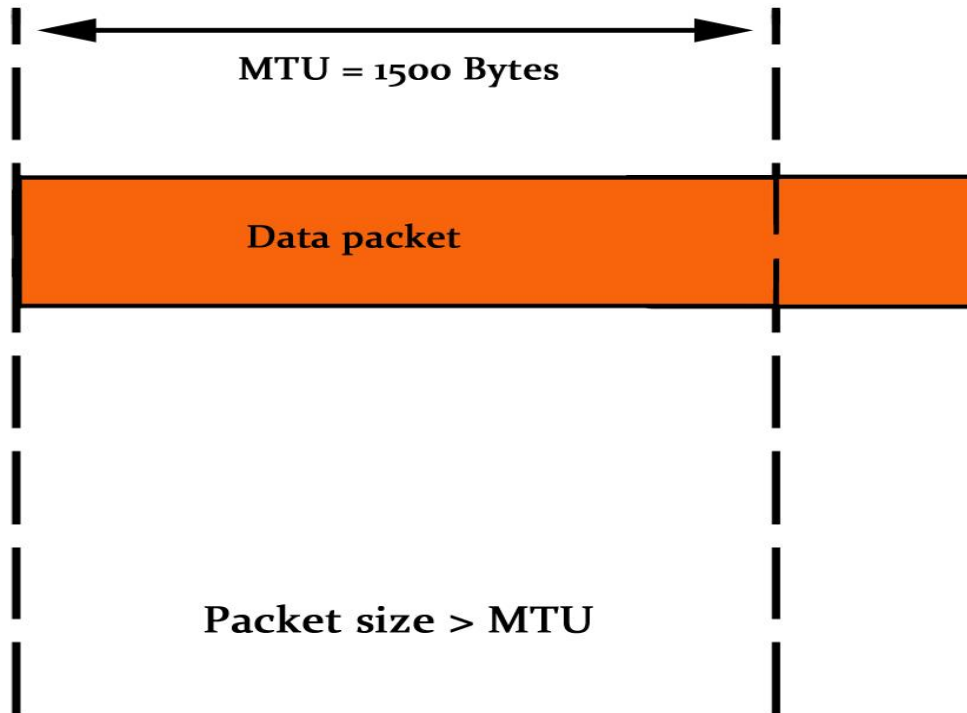
MTU : Maximum Transmission Unit
RTT: Round Trip Time

TIME – How it works



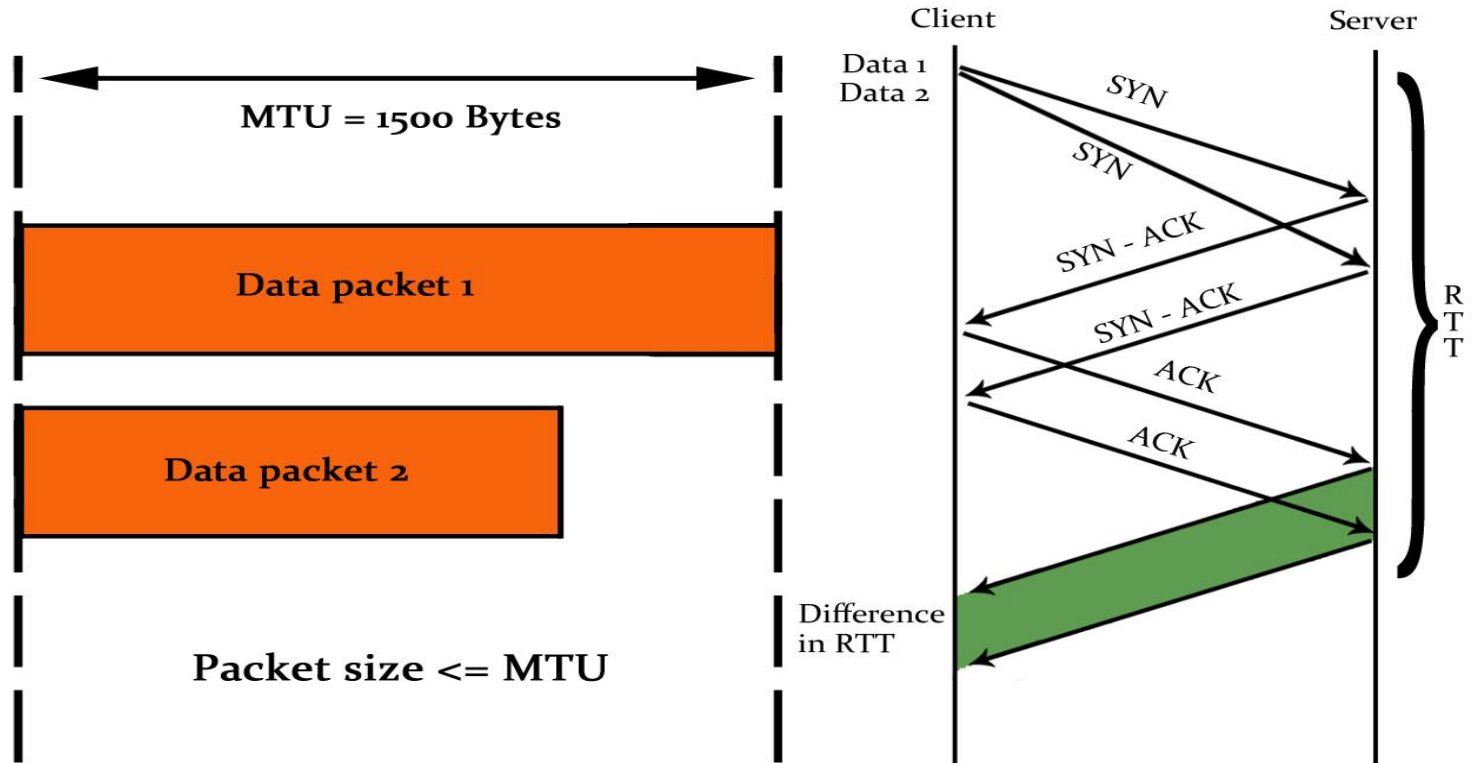
MTU : Maximum Transmission Unit
RTT: Round Trip Time

TIME – How it works



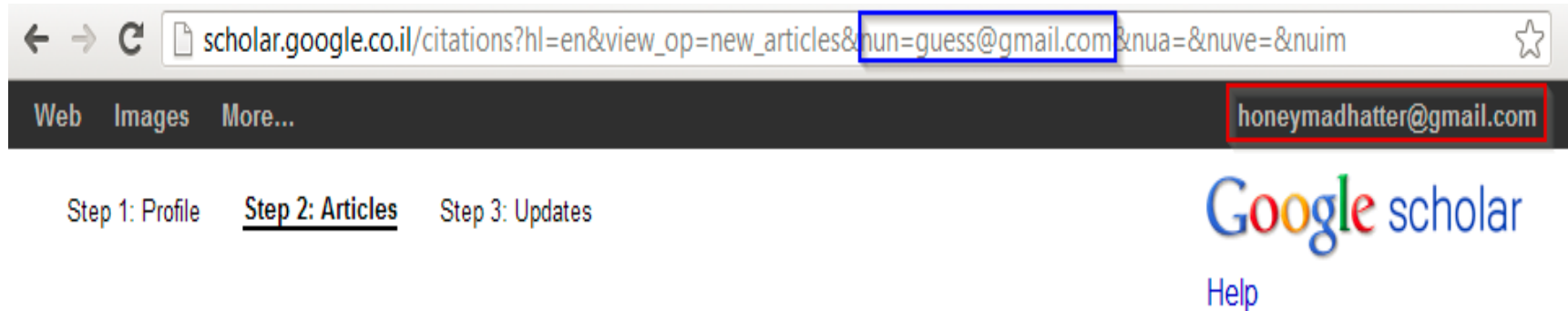
MTU : Maximum Transmission Unit
RTT: Round Trip Time

TIME – How it works



MTU : Maximum Transmission Unit
RTT: Round Trip Time

TIME – How it works



Add articles - guess@gmail.com

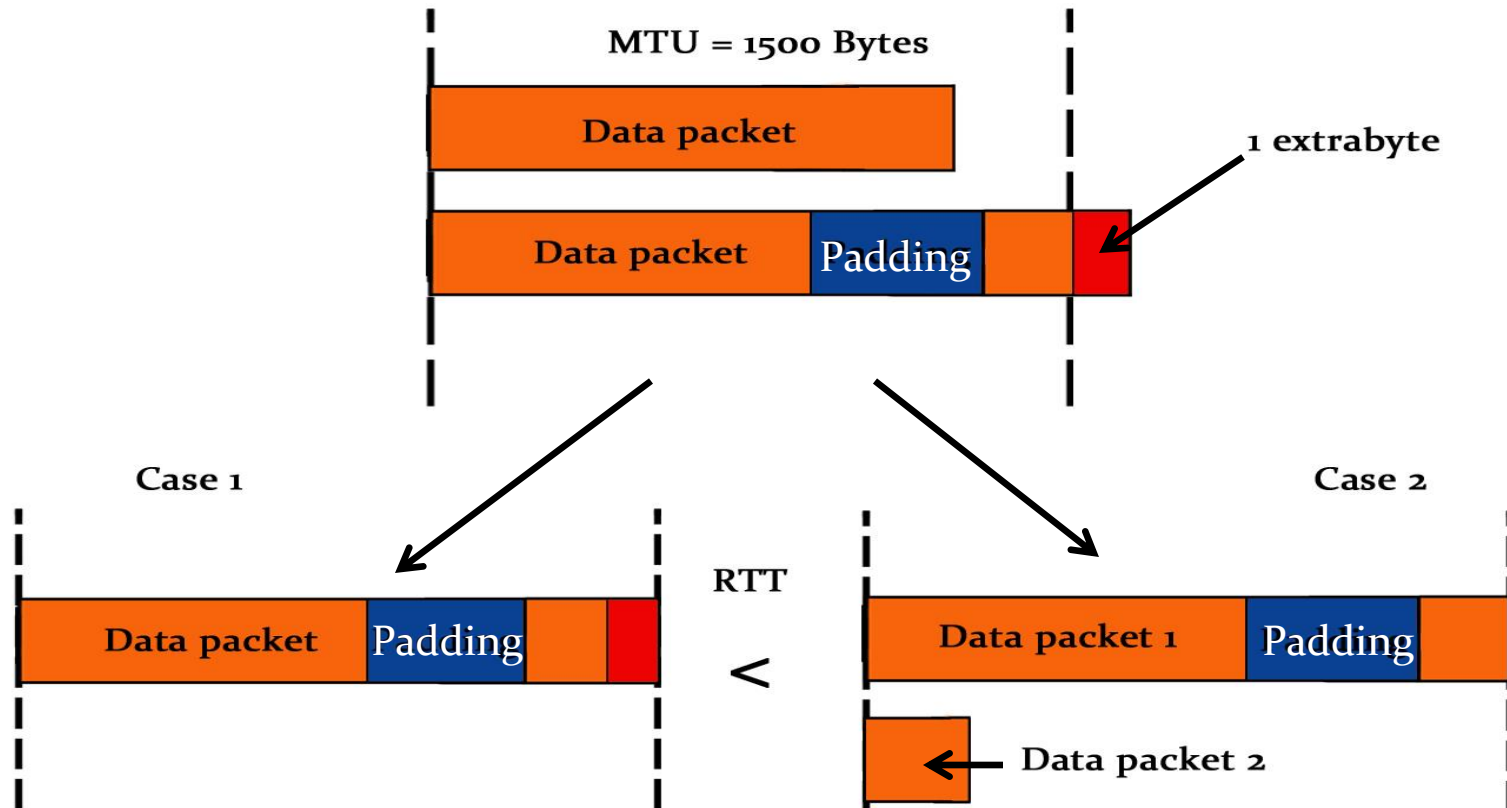
Find articles that you've written and add them to your profile. Later, you can edit or delete the articles in your profile or add more articles to your profile.

author: "guess@gmail.com"

Search article groups

Screenshot credit - Tal Be'ery BH presentation

TIME – How it works



MTU : Maximum Transmission Unit
RTT: Round Trip Time

TIME – Feasibility

- No requirements for Man-in-the-Middle
- Concentrate on HTTP responses
- The attacker creates HTTP request with JavaScript and response timing leaks the request size.
- Repeat for few times to void aberration due to network jitter.

TIME – Counter Measures

- Adding random timing delays to the decryption
- Browser should support and respect ``X-Frame-Options``
- Strict restriction on reflection of user input in the response.
- Enable anti-automation techniques like CAPTCHA, CSRF token

- Rate limiting using HAProxy
- Rate limiting via various DDOS protection

Comparison of counter measures

	BEAST	Lucky 13	RC4 Biases	CRIME	BREACH	TIME
CLIENT SIDE	Upgrade browsers			Upgrade browsers with no TLS compression support		Upgrade browsers with X-Frame-Options
SERVER SIDE		random timing delays	Throttle client initiated re-negotiations	Disable TLS compression	Mask the secret	Restrict reflection of user input
	Use RC4	Use RC4	Do not use RC4		anti-automation techniques	anti-automation techniques
	Upgrade to TLS 1.2	Upgrade to TLS 1.2	Upgrade to TLS 1.2		Separate secrets from user input	Random timing delay

Comparison of counter measures

	BEAST	Lucky 13	RC4 Biases	CRIME	BREACH	TIME
CLIENT SIDE	Upgrade browsers			Upgrade browsers with no TLS compression support		Upgrade browsers with X-Frame-Options
SERVER SIDE		random timing delays	Throttle client initiated re-negotiations	Disable TLS compression	Mask the secret	Restrict reflection of user input
	Use RC4	Use RC4	Do not use RC4		anti-automation techniques	anti-automation techniques
	Upgrade to TLS 1.2	Upgrade to TLS 1.2	Upgrade to TLS 1.2		Separate secrets from user input	Random timing delay

Comparison of counter measures

	BEAST	Lucky 13	RC ₄ Biases	CRIME	BREACH	TIME
CLIENT SIDE	Upgrade browsers			Upgrade browsers with no TLS compression support		Upgrade browsers with X-Frame-Options
SERVER SIDE			Throttle client initiated re-negotiations	Disable TLS compression	Mask the secret	Restrict reflection of user input
	Use RC ₄	Use RC ₄	Do not use RC ₄		anti-automation techniques	anti-automation techniques
	Upgrade to TLS 1.2	Upgrade to TLS 1.2	Upgrade to TLS 1.2			

Comparison of counter measures

	BEAST	Lucky 13	RC ₄ Biases	CRIME	BREACH	TIME
CLIENT SIDE	Upgrade browsers			Upgrade browsers with no TLS compression support		Upgrade browsers with X-Frame-Options
SERVER SIDE			Throttle client initiated re-negotiations	Disable TLS compression	Mask the secret	Restrict reflection of user input
	Use RC ₄	Use RC ₄	Do not use RC ₄		anti-automation techniques	anti-automation techniques
	Upgrade to TLS 1.2	Upgrade to TLS 1.2	Upgrade to TLS 1.2			

Comparison of counter measures

	BEAST	Lucky 13	RC4 Biases	CRIME	BREACH	TIME
CLIENT SIDE	Upgrade browsers			Upgrade browsers with no TLS compression support		Upgrade browsers with X-Frame-Options
SERVER SIDE			Throttle client initiated re-negotiations	Disable TLS compression	Mask the secret	Restrict reflection of user input
	Use RC4	Use RC4	Do not use RC4		anti-automation techniques	anti-automation techniques
	Upgrade to TLS 1.2	Upgrade to TLS 1.2	Upgrade to TLS 1.2			

Summarizing counter measures

	BEAST	Lucky 13	RC4 Biases	CRIME	BREACH	TIME
CLIENT SIDE	Upgrade browsers			Upgrade browsers with no TLS compression support		Upgrade browsers with X-Frame-Options
SERVER SIDE			Throttle client initiated re-negotiations	Disable TLS compression	Mask the secret	Restrict reflection of user input
	Use RC4	Use RC4	Do not use RC4		anti-automation techniques	anti-automation techniques
	Upgrade to TLS 1.1 +	Upgrade to TLS 1.2	Upgrade to TLS 1.2			

Thank You

- Special Thanks to Shawn, Tom, Michael, Javed, Jonathan, Tim, Josh, Alban, Ryan, Aaron and everybody in iSEC
- Whitepaper: Attacks on SSL (bit.ly/1cAqL7o)
- **Pratik Guha Sarkar**
 - Security Consultant at iSEC Partners
 - *psarkar@isecpartners.com | [@pragusa55](https://twitter.com/pragusa55)*
- **Shawn Fitzgerald**
 - Principal Security Consultant at iSEC Partners
 - *shawn@isecpartners.com*

